

The internal structure of natural numbers and one method for the definition of large prime numbers

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Abstract

It holds that every product of natural numbers can also be written as a sum. The inverse does not hold when 1 is excluded from the product. For this reason, the investigation of natural numbers should be done through their sum and not through their product. Such an investigation is presented in the present article. We prove that primes play the same role for odd numbers as the powers of 2 for even numbers, and vice versa. The following theorem is proven: "Every natural number, except for 0 and 1, can be uniquely written as a linear combination of consecutive powers of 2 with the coefficients of the linear combination being -1 or +1." This theorem reveals a set of symmetries in the internal order of natural numbers which cannot be derived when studying natural numbers on the basis of the product. From such a symmetry a method for identifying large prime numbers is derived.

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1. INTRODUCTION

It holds that every product of natural numbers can also be written as a sum. The inverse (i.e. each sum of natural numbers can be written as a product) does not hold when 1 is excluded from the product. This is due to prime numbers p which can be written as a product only in the form of $p = 1 \cdot p$. For this reason, the investigation of natural numbers should be done through their sum and not through their product. Such an investigation is presented in the present article.

We prove that each natural number can be written as a sum of three or more consecutive natural numbers except of the powers of 2 and the prime numbers. Each power of 2 and each prime number cannot be written as a sum of three or more consecutive natural numbers. Primes play the same role for odd numbers as the powers of 2 for even numbers, and vice versa.

We prove a theorem which is analogous to the fundamental theorem of arithmetic, when we study the positive integers with respect to addition: "Every natural number, with the exception of 0 and 1, can be written in a unique way as a linear combination of consecutive powers of 2, with the coefficients of the linear combination being -1 or +1." This theorem reveals a set of symmetries in the internal order of natural numbers which cannot be derived when studying natural numbers on the basis of the product. From such a symmetry a method for identifying large prime numbers is derived.

2. THE SEQUENCE $\mu(k,n)$

We consider the sequence of natural numbers

$$\begin{aligned}\mu(k,n) &= k + (k+1) + (k+2) + \dots + (k+n) = \frac{(n+1)(2k+n)}{2} \\ k \in \mathbb{N}^* &= \{1, 2, 3, \dots\} \\ n \in A &= \{2, 3, 4, \dots\}\end{aligned}. \quad (2.1)$$

For the sequence $\mu(k,n)$ the following theorem holds:

Theorem 2.1.

" For the sequence $\mu(k,n)$ the following hold:

1. $\mu(k,n) \in \mathbb{N}^*$.
2. No element of the sequence is a prime number.
3. No element of the sequence is a power of 2 .
4. The range of the sequence is all natural numbers that are not primes and are not powers of 2 .

Proof.

1. $\mu(k,n) \in \mathbb{N}^*$ as a sum of natural numbers.

2. $n \in A = \{2, 3, 4, \dots\}$ and therefore it holds that

$$n \geq 2$$

$$n+1 \geq 3$$

Also we have that

$$2k+n \geq 4$$

$$\frac{2k+n}{2} \geq \frac{3}{2} > 1$$

since $k \in \mathbb{N}^*$ and $n \in A = \{2, 3, 4, \dots\}$. Thus, the product

$$\frac{(n+1)(2k+n)}{2} = \mu(k,n)$$

is always a product of two natural numbers different than 1 , thus the natural number $\mu(k,n)$ cannot be prime.

3. Let that the natural number $\mu(k, n) = \frac{(n+1)(2k+n)}{2}$ is a power of 2 . Then, it exists

$\lambda \in \mathbb{N}$ such as

$$\begin{aligned} \frac{(n+1)(2k+n)}{2} &= 2^\lambda \\ (n+1)(2k+n) &= 2^{\lambda+1}. \end{aligned} \quad (2.2)$$

Equation (2.2) can hold if and only if there exist $\lambda_1, \lambda_2 \in \mathbb{N}$ such as

$$n+1 = 2^{\lambda_1} \wedge 2k+n = 2^{\lambda_2}$$

and equivalently

$$\left. \begin{array}{l} n = 2^{\lambda_1} - 1 \\ n = 2^{\lambda_2} - 2k \end{array} \right\}. \quad (2.3)$$

We eliminate n from equations (2.3) and we obtain

$$2^{\lambda_1} - 1 = 2^{\lambda_2} - 2k$$

and equivalently

$$2k - 1 = 2^{\lambda_2} - 2^{\lambda_1}$$

which is impossible since the first part of the equation is an odd number and the second part is an even number. Thus, the range of the sequence $\mu(k, n)$ does not include the powers of 2 .

4. We now prove that the range of the sequence $\mu(k, n)$ includes all natural numbers that are not primes and are not powers of 2 . Let a random natural number N which is not a prime nor a power of 2 . Then, N can be written in the form

$$N = \chi \psi$$

where at least one of the χ, ψ is an odd number ≥ 3 . Let χ be an odd number ≥ 3 . We will prove that there are always exist $k \in \mathbb{N}$ and $n \in A = \{2, 3, 4, \dots\}$ such as

$$N = \chi \cdot \psi = \mu(k, n).$$

We consider the following two pairs of k and n :

$$\chi \leq 2\psi - 1, \chi, \psi \in \mathbb{N}$$

$$k = k_1 = \frac{2\psi + 1 - \chi}{2} \quad (2.4)$$

$$n = n_1 = \chi - 1$$

$$\chi \geq 2\psi + 1, \chi, \psi \in \mathbb{N}$$

$$k = k_2 = \frac{\chi + 1 - 2\psi}{2} \quad (2.5)$$

$$n = n_2 = 2\psi - 1$$

For every $\chi, \psi \in \mathbb{N}$ it holds either the inequality $\chi \leq 2\psi - 1$ or the inequality $\chi \geq 2\psi + 1$. Thus, for each pair of naturals (χ, ψ) , where χ is odd, at least one of the pairs (k_1, n_1) , (k_2, n_2) of equations (2.4), (2.5) is defined. We now prove that “when the natural number k_1 of equation (2.4) is $k_1 = 0$ then the natural number k_2 of equation (2.5) is $k_2 = 1$ and additionally it holds that $n_2 > 2$ ”. For $k_1 = 0$ from equations (2.4) we take

$$\chi = 2\psi + 1$$

and from equations (2.5) we have that

$$k_2 = \frac{(2\psi + 1) + 1 - 2\psi}{2} = 1$$

$$n_2 = 2\psi - 1$$

and because $\psi \geq 2$ we obtain

$$k_2 = 1$$

$$n_2 = 2\psi - 1 \geq 3 > 2$$

We now prove that when $k_2 = 0$ in equations (2.5), then in equations (2.4) it is $k_1 = 1$ and $n_1 > 2$. For $k_2 = 0$, from equations (2.5) we obtain

$$\chi = 2\psi - 1$$

and from equations (2.4) we get

$$k_1 = \frac{2\psi + 1 - (2\psi - 1)}{2} = 1$$

$$n_1 = \chi - 1 = 2\psi - 2 \geq 2$$

We now prove that at least one of the k_1 and k_2 is positive. Let

$$k_1 < 0 \wedge k_2 < 0.$$

Then from equations (2.4) and (2.5) we have that

$$2\psi + 1 - \chi < 0 \wedge \chi + 1 - 2\psi < 0. \quad (2.6)$$

Taking into account that $\chi > 1$ is odd, that is $\chi = 2\rho + 1$, $\rho \in \mathbb{N}$, we obtain from inequalities (2.6)

$$2\psi + 1 - (2\rho - 1) < 0 \wedge (2\rho + 1) + 1 - 2\psi < 0$$

$$2\psi - 2\rho < 0 \wedge 2\rho - 2\psi + 2 > 0$$

$$\psi < \rho \wedge \psi > \rho + 1$$

which is absurd. Thus, at least one of k_1 and k_2 is positive.

For equations (2.4) we take

$$\begin{aligned}\mu(k_1, n_1) &= \frac{(n_1+1)(2k_1+n_1)}{2} \\ &= \frac{(\chi-1+1)\left(2\frac{2\psi+1-\chi}{2}+\chi-1\right)}{2} = \frac{\chi(2)\psi}{2} = \chi\psi = N\end{aligned}$$

For equations (2.5) we obtain

$$\begin{aligned}\mu(k_2, n_2) &= \frac{(n_2+1)(2k_2+n_2)}{2} \\ &= \frac{(2\psi-1+1)\left(2\frac{\chi+1-2\psi}{2}+2\psi-1\right)}{2} = \frac{2\psi\chi}{2} = \chi\psi = N\end{aligned}$$

Thus, there are always exist $k \in \mathbb{N}^*$ and $n \in A = \{2, 3, 4, \dots\}$ such as

$N = \chi\psi = \mu(k, n)$ for every N which is not a prime number and is not a power of 2. \square

Example 2.1. For the natural number $N = 40$ we have

$$N = 40 = 5 \cdot 8$$

$$\chi = 5$$

$$\psi = 8$$

and from equations (2.4) we get

$$k = k_1 = \frac{16+1-5}{2} = 6$$

$$n = n_1 = 5 - 1 = 4$$

thus, we obtain

$$40 = \mu(6, 4).$$

Example 2.2. For the natural number $N = 51$,

$$N = 51 = 3 \cdot 17 = 17 \cdot 3$$

there are two cases. First case:

$$N = 51 = 3 \cdot 17$$

$$\chi = 3$$

$$\psi = 17$$

and from equations (2.4) we obtain

$$k = k_1 = \frac{34+1-3}{2} = 16$$

$$n = n_1 = 3 - 1 = 2$$

thus,

$$51 = \mu(16, 2).$$

Second case:

$$N = 51 = 17 \cdot 3$$

$$\chi = 17$$

$$\psi = 3$$

and from equations (2.5) we obtain

$$k = k_2 = \frac{17+1-6}{2} = 6$$

$$n = n_2 = 6 - 1 = 5$$

thus,

$$51 = \mu(6,5).$$

The second example expresses a general property of the sequence $\mu(k,n)$. The more composite an odd number that is not prime (or an even number that is not a power of 2) is, the more are the $\mu(k,n)$ combinations that generate it.

Example 2.3.

$$135 = 15 \cdot 9 = 27 \cdot 5 = 9 \cdot 15 = 45 \cdot 3 = 5 \cdot 27 = 3 \cdot 45$$

$$135 = \mu(2,14) = \mu(9,9) = \mu(11,8) = \mu(20,5) = \mu(25,4) = \mu(44,2)$$

$$\text{a. } 135 = 9 \cdot 15 = \mu(2,14) = \mu(11,8)$$

$$135 = 2 + 3 + 4 + \dots + 15 + 16 = 11 + 12 + 13 + \dots + 18 + 19.$$

$$\text{b. } 135 = 5 \cdot 27 = \mu(9,9) = \mu(25,4)$$

$$135 = 9 + 10 + 11 + \dots + 17 + 18 = 25 + 26 + 27 + 28 + 29.$$

$$\text{c. } 135 = 3 \cdot 45 = \mu(20,5) = \mu(44,2)$$

$$135 = 20 + 21 + 22 + 23 + 24 + 25 = 44 + 45 + 46.$$

In the transitive property of multiplication, when writing a composite odd number or an even number that is not a power of 2 as a product of two natural numbers, we use the same natural numbers $\chi, \psi \in \mathbb{N}$:

$$\Phi = \chi \cdot \psi = \psi \cdot \chi.$$

On the contrary, the natural number Φ can be written in the form $\Phi = \mu(k,n)$ using different natural numbers $k \in \mathbb{N}^*$ and $n \in A = \{2, 3, 4, \dots\}$, through equations (2.4), (2.5). This difference between the product and the sum can also become evident in example 2.3:

$$135 = 3 \cdot 45 = 45 \cdot 3$$

$$135 = 44 + 45 + 46 = 20 + 21 + 22 + 23 + 24 + 25$$

From Theorem 2.1 the following corollary is derived:

Corollary 2.1. “1. Every natural number which is not a power of 2 and is not a prime can be written as the sum of three or more *consecutive* natural numbers.

2. Every power of 2 and every prime number cannot be written as the sum of three or more *consecutive* natural numbers.”

Proof. Corollary 2.1 is a direct consequence of Theorem 2.1. \square

3. THE CONCEPT OF REARRANGEMENT

In this paragraph, we present the concept of rearrangement of the composite odd numbers and even numbers that are not power of 2. Moreover, we prove some of the consequences of the rearrangement in the Diophantine analysis. The concept of rearrangement is given from the following definition:

Definition. “We say that the sequence $\mu(k, n), k \in \mathbb{N}^*, n \in A = \{2, 3, 4, \dots\}$ is rearranged if there exist natural numbers $k_1 \in \mathbb{N}^*, n_1 \in A, (k_1, n_1) \neq (k, n)$ such as

$$\mu(k, n) = \mu(k_1, n_1). \quad (3.1)$$

From equation (2.1) written in the form of

$$\mu(k, n) = k + (k+1) + (k+2) + \dots + (k+n)$$

two different types of rearrangement are derived: The “compression”, during which n decreases with a simultaneous increase of k . The «decompression», during which n increases with a simultaneous decrease of k . The following theorem provides the criterion for the rearrangement of the sequence $\mu(k, n)$.

Theorem 3.1. “’1. The sequence $\mu(k_1, n_1), (k_1, n_1) \in \mathbb{N}^* \times A$ can be compressed

$$\mu(k_1, n_1) = \mu(k_1 + \varphi, n_1 - \omega) \quad (3.2)$$

if and only if there exist $\varphi, \omega \in \mathbb{N}^*, \omega \leq n_1 - 2$ which satisfies the equation

$$\begin{aligned} \omega^2 - (2k_1 + 2n_1 + 1 + 2\varphi)\omega + 2(n_1 + 1)\varphi &= 0 \\ \varphi, \omega \in \mathbb{N}^* &. \\ \omega &\leq n_1 - 2 \end{aligned} \quad (3.3)$$

2. The sequence $\mu(k_2, n_2), (k_2, n_2) \in \mathbb{N}^* \times A$ can be decompressed

$$\mu(k_2, n_2) = \mu(k_2 - \varphi, n_2 + \omega) \quad (3.4)$$

if and only if there exist $\varphi, \omega \in \mathbb{N}^*, \varphi \leq k_2 - 1$ which satisfies the equation

$$\begin{aligned} \omega^2 + (2k_2 + 2n_2 + 1 - 2\varphi)\omega - 2(n_2 + 1)\varphi &= 0 \\ \varphi, \omega \in \mathbb{N}^* \\ \varphi \leq k_2 - 1 \end{aligned} . \quad (3.5)$$

3. The odd number $\Pi \neq 1$ is prime if and only if the sequence

$$\begin{aligned} \mu(k, n) &= \Pi \cdot 2^l \\ l, k \in \mathbb{N}^*, n \in A \end{aligned} \quad (3.6)$$

cannot be rearranged.

4. The odd Π is prime if and only if the sequence

$$\mu\left(\frac{\Pi+1}{2}, \Pi-1\right) = \Pi^2 \quad (3.7)$$

cannot be rearranged."

Proof. 1,2. We prove part 1 of the corollary and similarly number 2 can also be proven. From equation (4.1) we conclude that the sequence $\mu(k_1, n_1)$ can be compressed if and only if there exist $\varphi, \omega \in \mathbb{N}^*$ such as

$$\mu(k_1, n_1) = \mu(k_1 + \varphi, n_1 - \omega).$$

In this equation the natural number $n_1 - \omega$ belongs to the set $A = \{2, 3, 4, \dots\}$ and thus $n_1 - \omega \geq 2 \Leftrightarrow \omega \leq n_1 - 2$. Next, from equations (2.1) we obtain

$$\begin{aligned} \mu(k_1, n_1) &= \mu(k_1 + \varphi, n_1 - \omega) \\ \frac{(n_1 + 1)(2k_1 + n_1)}{2} &= \frac{(n_1 - \omega + 1)[2(k_1 + \varphi) + n_1 - \omega]}{2} \end{aligned}$$

and after the calculations we get equation (3.3).

3. The sequence (3.6) is derived from equations (2.4) or (2.5) for $\chi = \Pi$ and $\psi = 2^l$. Thus, in the product $\chi\psi$ the only odd number is Π . If the sequence $\mu(k, n)$ in equation (3.6) cannot be rearranged then the odd number Π has no divisors. Thus, Π is prime. Obviously, the inverse also holds.

4. First, we prove equations (3.7). From equation (2.1) we obtain:

$$\mu\left(\frac{\Pi+1}{2}, \Pi-1\right) = \frac{(\Pi-1+1)\left(2\frac{\Pi+1}{2} + \Pi - 1\right)}{2} = \Pi^2.$$

In case that the odd number Π is prime in equations (2.4), (2.5) the natural numbers χ, ψ are unique

$\chi = \Pi \wedge \psi = \Pi$, and from equation (2.5) we get $k = \frac{\Pi+1}{2} \wedge n = \Pi - 1$. Thus, the sequence

$\mu(k, n) = \mu\left(\frac{\Pi+1}{2}, \Pi-1\right)$ cannot be rearranged. Conversely, if the sequence

$\mu\left(\frac{\Pi+1}{2}, \Pi-1\right) = \Pi^2 = \Pi \cdot \Pi$ cannot be rearranged the odd number Π cannot be composite and

thus Π is prime. \square

We now prove the following corollary:

Corollary 3.1. "1. The odd number Φ ,

$$\begin{aligned} \Phi &= \Pi^2 = \mu\left(\frac{\Pi+1}{2}, \Pi-1\right) \\ \Pi &= \text{odd} \\ \Pi &\neq 1 \end{aligned} \tag{3.8}$$

is decompressed and compressed if and only if the odd number Π is composite.

2. The even number α_1 ,

$$\begin{aligned} \alpha_1 &= 2^l \Pi = \mu\left(2^l - \frac{\Pi-1}{2}, \Pi-1\right) \\ \Pi &= \text{odd} \\ 3 \leq \Pi &\leq 2^l - 1 \\ l \in \mathbb{N}, l &\geq 2 \end{aligned} \tag{3.9}$$

cannot be decompressed, while it compresses if and only if the odd number Π is composite.

3. The even number α_2 ,

$$\begin{aligned} \alpha_2 &= 2^l \Pi = \mu\left(\frac{\Pi+1}{2} - 2^l, 2^{l+1} - 1\right) \\ \Pi &= \text{odd} \\ \Pi &\geq 2^{l+1} + 1 \\ l \in \mathbb{N}^* & \end{aligned} \tag{3.10}$$

cannot be compressed, while it decompresses if and only if the odd number Π is composite.

4. Every even number that is not a power of can be written either in the form of equation (3.9) or in the form of equation (3.10)."

Proof.

1. It is derived directly through number (4) of Theorem 3.1. A second proof can be derived through equations (2.4), (2.5) since every composite odd Π can be written in the form of $\Pi = \chi\psi$, $\chi, \psi \in \mathbb{N}$, χ, ψ odds.

2,3.

Let the even number α ,

$$\begin{aligned}\alpha &= 2^l \Pi \\ \Pi &= \text{odd} . \\ l &\in \mathbb{N}^*\end{aligned}\tag{3.11}$$

From equation (2.4) we obtain

$$\begin{aligned}k &= \frac{2 \cdot 2^l + 1 - \Pi}{2} = 2^l - \frac{\Pi - 1}{2} \\ n &= \Pi - 1\end{aligned}\tag{3.12}$$

and since $k, n \in \mathbb{N}, k \geq 1 \wedge n \geq 2$ we get

$$\begin{aligned}\frac{2 \cdot 2^l + 1 - \Pi}{2} &\geq 1 \\ \Pi - 1 &\geq 2\end{aligned}$$

and equivalently

$$3 \leq \Pi \leq 2^{l+1} - 1.$$

In the second of equations (3.12) the natural number n obtains the maximum possible value of $n = \Pi - 1$, and thus the natural number k takes the minimum possible value in the first of equations (3.12). Thus, the even number

$$\alpha_1 = \mu\left(2^l - \frac{\Pi - 1}{2}, \Pi - 1\right)$$

cannot decompress. If the odd number Π is composite then it can be written in the form of $\Pi = \chi\psi$, $\chi, \psi \in \mathbb{N}^*$, χ, ψ odds, $\chi, \psi < \Pi$, $\alpha_1 = 2^l \chi\psi$. Therefore, the natural number $\alpha_1 = 2^l \chi\psi$ decompresses since from equations (3.11) it can be written in the form of $\alpha_1 = \mu(k, n)$ with $n = \chi - 1 < \Pi - 1$. Similarly, the proof of 3 is derived from equations (2.5).

4. From the above proof process it follows that every even number that is not a power of 2 can be written either in the form of equation (3.9) or in the form of equation (3.10). \square

By substituting $\Pi = P = \text{prime}$ in equations of Theorem 3.1 and of corollary 3.1 four sets of equations are derived, each including infinite *impossible* diophantine equations.

Example 3.1. The odd number $P = 999961$ is prime. Thus, combining (1) of Theorem 3.1 with (1) of corollary 3.1 we conclude that there is no pair $(\omega, \varphi) \in \mathbb{N}^2$ with $\omega \leq 999958$ which satisfies the diophantine equation

$$\omega^2 - (2999883 + 2\varphi)\omega + 1999922\varphi = 0.$$

We now prove the following corollary:

Corollary 3.2 "The square of every prime number can be uniquely written as the sum of consecutive natural numbers."

Proof. For $\Pi = P = \text{prime}$ in equation (3.5) we obtain

$$P^2 = \mu\left(\frac{P+1}{2}, P-1\right). \quad (3.13)$$

According with 4 of Theorem 3.1 the odd P^2 cannot be rearranged. Thus, the odd can be uniquely written as the sum of consecutive natural numbers, as given from equation (3.13). \square

Example 3.2. The odd $P = 17$ is prime. From equation (3.13) for $P = 17$ we obtain

$$289 = \mu(9, 16)$$

and from equation (2.1) we get

$$289 = 9 + 10 + 11 + 12 + 13 + 14 + 15 + 16 + 17 + 18 + 19 + 20 + 21 + 22 + 23 + 24 + 25$$

which is the only way in which the odd number 289 can be written as a sum of consecutive natural numbers.

4. NATURAL NUMBERS AS LINEAR COMBINATION OF CONSECUTIVE POWERS OF 2

According to the fundamental theorem of arithmetic, every natural number can be uniquely written as a product of powers of prime numbers. The previously presented study reveals a correspondence between odd prime numbers and the powers of 2. Thus, the question arises whether there exists a theorem for the powers of 2 corresponding to the fundamental theorem of arithmetic. The answer is given by the following theorem:

Theorem 4.1. 'Every natural number, with the exception of 0 and 1, can be uniquely written as a linear combination of consecutive powers of 2, with the coefficients of the linear combination being -1 or +1.'

Proof. Let the odd number Π as given from equation

$$\Pi = \Pi(\nu, \beta_i) = 2^{\nu+1} + 2^\nu \pm 2^{\nu-1} \pm 2^{\nu-2} \pm \dots \pm 2^1 \pm 2^0 = 2^{\nu+1} + 2^\nu + \sum_{i=0}^{\nu-1} \beta_i 2^i$$

$$\beta_i = \pm 1, i = 0, 1, 2, \dots, \nu - 1$$

$$\nu \in \mathbb{N}$$
(4.1)

From equation (4.1) for $\nu = 0$ we obtain

$$\Pi = 2^1 + 2^0 = 2 + 1 = 3.$$

We now examine the case where $\nu \in \mathbb{N}^*$. The lowest value that the odd number Π of equation (4.1) can obtain is

$$\begin{aligned}\Pi_{\min} &= \Pi(\nu) = 2^{\nu+1} + 2^\nu - 2^{\nu-1} - 2^{\nu-2} - \dots - 2^1 - 1 \\ \Pi_{\min} &= \Pi(\nu) = 2^{\nu+1} + 1.\end{aligned}\tag{4.2}$$

The largest value that the odd number Π of equation (4.1) can obtain is

$$\begin{aligned}\Pi_{\max} &= \Pi(\nu) = 2^{\nu+1} + 2^\nu + 2^{\nu-1} + \dots + 2^1 + 1 \\ \Pi_{\max} &= \Pi(\nu) = 2^{\nu+2} - 1.\end{aligned}\tag{4.3}$$

Thus, for the odd numbers $\Pi = \Pi(\nu, \beta_i)$ of equation (4.1) the following inequality holds

$$\Pi_{\min} = 2^{\nu+1} + 1 \leq \Pi(\nu, \beta_i) \leq 2^{\nu+2} - 1 = \Pi_{\max}. \tag{4.4}$$

The number $N(\Pi(\nu, \beta_i))$ of odd numbers in the closed interval $[2^{\nu+1} + 1, 2^{\nu+2} - 1]$ is

$$\begin{aligned}N(\Pi(\nu, \beta_i)) &= \frac{\Pi_{\max} - \Pi_{\min}}{2} + 1 = \frac{(2^{\nu+2} - 1) - (2^{\nu+1} + 1)}{2} + 1 \\ N(\Pi(\nu, \beta_i)) &= 2^\nu.\end{aligned}\tag{4.5}$$

The integers $\beta_i, i = 0, 1, 2, \dots, \nu - 1$ in equation (4.1) can take only two values, $\beta_i = -1 \vee \beta_i = +1$, thus equation (4.1) gives exactly $2^\nu = N(\Pi(\nu, \beta_i))$ odd numbers. Therefore, for every $\nu \in \mathbb{N}^*$ equation (4.1) gives all odd numbers in the interval $[2^{\nu+1} + 1, 2^{\nu+2} - 1]$.

We now prove the theorem for the even numbers. Every even number α which is a power of 2 can be uniquely written in the form of $\alpha = 2^\nu, \nu \in \mathbb{N}^*$. We now consider the case where the even number α is not a power of 2. In that case, according to corollary 3.1 the even number α is written in the form of

$$\alpha = 2^l \Pi, \Pi = \text{odd}, \Pi \neq 1, l \in \mathbb{N}^*. \tag{4.6}$$

We now prove that the even number α can be uniquely written in the form of equation (4.6). If we assume that the even number α can be written in the form of

$$\begin{aligned} \alpha &= 2^l \Pi = 2^{l'} \Pi' \\ l &\neq l' (l > l') \\ \Pi &\neq \Pi' \\ l, l' &\in \mathbb{N}^* \\ \Pi, \Pi' &= \text{odd} \end{aligned} \tag{4.7}$$

then we obtain

$$2^l \Pi = 2^{l'} \Pi'$$

$$2^{l-l'} \Pi = \Pi'$$

which is impossible, since the first part of this equation is even and the second odd. Thus, it is $l = l'$ and we take that $\Pi = \Pi'$ from equation (4.7). Therefore, every even number α that is not a power of 2 can be uniquely written in the form of equation (4.6). The odd number Π of equation (4.6) can be uniquely written in the form of equation (4.1), thus from equation (4.6) it is derived that every even number α that is not a power of 2 can be uniquely written in the form of equation

$$\begin{aligned} \alpha &= \alpha(l, v, \beta_i) = 2^l \left(2^{v+1} + 2^v + \sum_{i=0}^{v-1} \beta_i 2^i \right) \\ l &\in \mathbb{N}^*, v \in \mathbb{N} \\ \beta_i &= \pm 1, i = 0, 1, 2, \dots, v-1 \end{aligned} \tag{4.8}$$

and equivalently

$$\begin{aligned} \alpha &= \alpha(l, v, \beta_i) = 2^{l+v+1} + 2^{l+v} + \sum_{i=0}^{v-1} \beta_i 2^{l+i} \\ l &\in \mathbb{N}^*, v \in \mathbb{N} \\ \beta_i &= \pm 1, i = 0, 1, 2, \dots, v-1 \end{aligned} \tag{4.9}$$

For 1 we take

$$1 = 2^0$$

$$1 = 2^1 - 2^0$$

thus, it can be written in two ways in the form of equation (4.1). Both the odds of equation (4.1) and the evens of the equation (4.8) are positive. Thus, 0 cannot be written either in the form of equation (4.1) or in the form of equation (4.8). \square

In order to write an odd number $\Pi \neq 1, 3$ in the form of equation (4.1) we initially define the $v \in \mathbb{N}^*$ from inequality (4.4). Then, we calculate the sum

$$2^{\nu+1} + 2^\nu.$$

If it holds that $2^{\nu+1} + 2^\nu < \Pi$ we add the $2^{\nu-1}$, whereas if it holds that $2^{\nu+1} + 2^\nu > \Pi$ then we subtract it. By repeating the process exactly ν times we write the odd number Π in the form of equation (4.1). The number of ν steps needed in order to write the odd number Π in the form of equation (4.1) is extremely low compared to the magnitude of the odd number Π , as derived from inequality (4.4).

Example 4.1. For the odd number $\Pi = 23$ we obtain from inequality (4.4)

$$2^{\nu+1} + 1 < 23 < 2^{\nu+2} - 1$$

$$2^{\nu+1} + 2 < 24 < 2^{\nu+2}$$

$$2^\nu < 12 < 2^{\nu+1}$$

thus $\nu = 3$. Then, we have

$$2^{\nu+1} + 2^\nu = 2^4 + 2^3 = 24 > 23 \text{ (thus } 2^2 \text{ is subtracted)}$$

$$2^4 + 2^3 - 2^2 = 20 < 23 \text{ (thus } 2^1 \text{ is added)}$$

$$2^4 + 2^3 - 2^2 + 2^1 = 22 < 23 \text{ (thus } 2^0 = 1 \text{ is added)}$$

$$2^4 + 2^3 - 2^2 + 2^1 + 1 = 23.$$

Fermat numbers F_s can be written directly in the form of equation (4.1), since they are of the form Π_{\min} ,

$$F_s = 2^{2^s} + 1 = \Pi_{\min} (2^s - 1) = 2^{2^s} + 2^{2^s-1} - 2^{2^s-2} - 2^{2^s-3} - \dots - 2^1 - 1. \quad (4.10)$$

$s \in \mathbb{N}^*$

Mersenne numbers M_p can be written directly in the form of equation (4.1), since they are of the form Π_{\max} ,

$$M_p = 2^p - 1 = \Pi_{\max} (p - 2) = 2^{p-1} + 2^{p-2} + 2^{p-3} + \dots + 2^1 + 1. \quad (4.11)$$

$p = \text{prime}$

In order to write an even number α that is not a power of 2 in the form of equation (4.1), initially it is consecutively divided by 2 and it takes of the form of equation (4.6). Then, we write the odd number Π in the form of equation (4.1).

Example 4.2. By consecutively dividing the even number $\alpha = 368$ by 2 we obtain $\alpha = 368 = 2^4 \cdot 23$. Then, we write the odd number $\Pi = 23$ in the form of equation (4.1), $23 = 2^4 + 2^3 - 2^2 + 2^1 + 1$, and we get

$$368 = 2^4 (2^4 + 2^3 - 2^2 + 2^1 + 1).$$

$$368 = 2^8 + 2^7 - 2^6 + 2^5 + 2^4$$

This equation gives the unique way in which the even number $\alpha = 368$ can be written in the form of equation (4.9).

From inequality (4.4) we obtain

$$2^{\nu+1} + 1 \leq \Pi \leq 2^{\nu+2} - 1$$

$$2^{\nu+1} < 2^{\nu+1} + 1 \leq \Pi \leq 2^{\nu+2} - 1 < 2^{\nu+2}$$

$$2^{\nu+1} < \Pi < 2^{\nu+2}$$

$$(\nu + 1) \log 2 < \log \Pi < (\nu + 2) \log 2$$

from which we get

$$\frac{\log \Pi}{\log 2} - 1 < \nu + 1 < \frac{\log \Pi}{\log 2}$$

and finally

$$\nu + 1 = \left[\frac{\log \Pi}{\log 2} \right] \quad (4.12)$$

'where $\left[\frac{\log \Pi}{\log 2} \right]$ the integer part of $\frac{\log \Pi}{\log 2} \in \mathbb{R}$.

We now give the following definition:

Definition 4.1. (The fundamental Π^* symmetry) We define as the conjugate of the odd

$$\begin{aligned} \Pi = \Pi(\nu, \beta_i) &= 2^{\nu+1} + 2^\nu + \sum_{i=0}^{\nu-1} \beta_i 2^i \\ \beta_i &= \pm 1, i = 0, 1, 2, \dots, \nu - 1 \\ \nu &\in \mathbb{N}^* \end{aligned} \quad (4.13)$$

the odd Π^* ,

$$\begin{aligned} \Pi^* = \Pi^*(\nu, \gamma_j) &= 2^{\nu+1} + 2^\nu + \sum_{j=0}^{\nu-1} \gamma_j 2^j \\ \gamma_j &= \pm 1, j = 0, 1, 2, \dots, \nu - 1 \\ \nu &\in \mathbb{N}^* \end{aligned} \quad (4.14)$$

for which it holds

$$\gamma_k = -\beta_k \forall k = 0, 1, 2, \dots, \nu - 1. \quad (4.15)$$

For conjugate odds, the following corollary holds:

Corollary 4.1. " For the conjugate odds $\Pi = \Pi(\nu, \beta_i)$ and $\Pi^* = \Pi^*(\nu, \gamma_i)$ the following hold:

$$1. (\Pi^*)^* = \Pi . \quad (4.16)$$

$$2. \Pi + \Pi^* = 3 \cdot 2^{\nu+1} . \quad (4.17)$$

3. Π is divisible by 3 if and only if Π^* is divisible by 3."

Proof. 1.The 1 of the corollary is an immediate consequence of definition 4.1.

2. From equations (4.13), (4.14) and (4.15) we get

$$\Pi + \Pi^* = (2^{\nu+1} + 2^\nu) + (2^{\nu+1} + 2^\nu)$$

and, equivalently

$$\Pi + \Pi^* = 3 \cdot 2^{\nu+1} .$$

3. If the odd Π is divisible by 3 then it is written in the form $\Pi = 3x, x = \text{odd}$ and from equation (4.17) we get $3x + \Pi^* = 3 \cdot 2^{\nu+1}$ and equivalently $\Pi^* = 3(2^{\nu+1} - x)$. Similarly we can prove the inverse. \square

5. THE HARMONIC ODD NUMBERS AND A METHOD FOR DEFINING LARGE PRIME NUMBERS

The harmonic symmetry: We define as harmonic the odd numbers of equation (4.1) for which the signs of $\beta_i = \pm 1, i = 0, 1, 2, 3, \dots, \nu-1$ alternate:

$$\begin{aligned} \Pi_1 &= 2^{\nu+1} + 2^\nu - 2^{\nu-1} + 2^{\nu-2} - \dots - 2^1 + 1 = \frac{2^{\nu+3} + 1}{3} \\ \Pi_2 &= 2^{\nu+1} + 2^\nu + 2^{\nu-1} - 2^{\nu-2} + \dots + 2^1 - 1 = \frac{5 \times 2^{\nu+1} - 1}{3} . \end{aligned} \quad (5.1)$$

$$\nu = 2\lambda, \lambda \in \mathbb{N}^*$$

$$\begin{aligned} \Pi_1 &= 2^{\nu+1} + 2^\nu - 2^{\nu-1} + 2^{\nu-2} - \dots + 2^1 - 1 \\ \Pi_2 &= 2^{\nu+1} + 2^\nu + 2^{\nu-1} - 2^{\nu-2} + \dots - 2^1 + 1 . \end{aligned} \quad (5.2)$$

$$\nu = 2\lambda + 1, \lambda \in \mathbb{N}^*$$

From equations (5.1), (5.2) and definition 4.1 we obtain

$$\Pi_2 = \Pi_1^* = 3 \times 2^{\nu+1} - \Pi_1 \quad (5.3)$$

for the pair of harmonic odd numbers.

A method for the determination of large prime numbers emerges from the study we presented. This method is completely different from previous methods [1-5].When we consider the prime factorization of the odd integers

$$\Phi_1 = 2 + \Pi_1 = 2 + \frac{2^{\nu+3} + 1}{3} \quad (5.4)$$

$\nu = 2\lambda, \lambda \in \mathbb{N}^*$

$$\Phi_2 = -2 + \Pi_2 = -2 + \frac{5 \times 2^{\nu+1} - 1}{3} \quad (5.5)$$

$\nu = 2\lambda, \lambda \in \mathbb{N}^*$

$$\Phi_2 = \Phi_1^* = 3 \times 2^{\nu+1} - \Phi_1 \quad (5.6)$$

we have the following statement:

The factors of either Φ_1 or $\Phi_2 = \Phi_1^*$ consist of a set of small prime factors and one large factor. Hence from the factorization of Φ_1 and $\Phi_2 = \Phi_1^*$ of equations (5.4), (5.5) we get a large prime number.

Following are 11 examples where we have chosen arbitrary even ν , $600 \leq \nu \leq 1000$, in equations (5.4), (5.5).

1. $\nu = 604$

$$\begin{aligned} \Phi_1 = & 3 \times 5 \times 1423 \times 2677 \times 1039667 \times 1465469 \times 2033624 \times 136455 \times 907062 \times 140355 \times 606581 \times 460617 \\ & 960329 \times 378244 \times 909713 \times 340374 \times 546035 \times 722007 \times 834481 \times 807880 \times 893223 \times 943637 \times 129816 \times 307143 \times 853799 \\ & 454110 \times 280390 \times 103176 \times 523414 \times 883666 \times 509589 \times 711687 \times 765791. \end{aligned}$$

2. $\nu = 626$

$$\begin{aligned} \Phi_1 = & 13 \times 186653 \times 306 \times 032599 \times 340492 \times 581270 \times 323029 \times 570138 \times 222136 \times 733600 \times 420877 \times 092183 \times 139417 \\ & 574185 \times 782109 \times 955578 \times 496315 \times 765962 \times 131603 \times 014089 \times 519221 \times 871827 \times 181120 \times 845674 \times 859725 \times 387186 \\ & 219442 \times 305406 \times 755275 \times 821605 \times 426602 \times 403741 \times 599957. \end{aligned}$$

3. $\nu = 644$

$$\begin{aligned} \Phi_1 = & 5 \times 79 \times 12 \times 671297 \times 38892 \times 671359 \times 559494 \times 324882 \times 180204 \times 888273 \times 078001 \times 950134 \times 412751 \\ & 881230 \times 225550 \times 378061 \times 442396 \times 379471 \times 711953 \times 850899 \times 474720 \times 409489 \times 565536 \times 036909 \times 109253 \times 945965 \\ & 590266 \times 361910 \times 559333 \times 142120 \times 493266 \times 182138 \times 997818 \times 136400 \times 630503. \end{aligned}$$

4. $\nu = 688$

$$\begin{aligned} \Phi_1 = & 3^3 \times 5 \times 137 \times 2357 \times 84239 \times 14 \times 276659 \times 111598 \times 463167 \times 164995 \times 567141 \times 3547 \times 493034 \\ & 864246 \times 374604 \times 223939 \times 439254 \times 117526 \times 195183 \times 644765 \times 258504 \times 745395 \times 441461 \times 348003 \times 624541 \times 265182 \\ & 053620 \times 319595 \times 210678 \times 493117 \times 621150 \times 188802 \times 864705 \times 030169 \times 622562 \times 000148 \times 389984 \times 593085 \times 080457. \end{aligned}$$

5. $\nu = 732$

$$\begin{aligned} \Phi_1 = & 5^5 \times 19 \times 4357 \times 10093 \times 2 \times 901193 \times 373 \times 058471 \times 21 \times 318693 \times 003272 \times 810610 \times 223875 \times 009176 \\ & 985967 \times 454565 \times 131359 \times 547239 \times 807330 \times 702392 \times 207730 \times 665072 \times 351378 \times 572215 \times 387223 \times 133953 \times 567092 \end{aligned}$$

456647 869354 874941 347502 746701 928543 247909 407783 975122 056127 018272 539991 430427
637981.

6. $\nu = 818$

$\Phi_2 = 5 826599 918309 521729 414628 892756 111346 582385 085483 938095 996388 692690 239258$
 $901551 139189 714409 550909 093308 382061 608683 211278 156913 402724 889465 422572 029940$
 $710372 011896 628413 739441 379150 647273 555252 328499 350633 086191 299627 798178 454098$
 $229036 211569 513811$ is prime.

7. $\nu = 838 ?$

8. $\nu = 842$

$\Phi_1 = 13 \times 811 \times 7789 \times 15271 \times 66809 \times 933 419184 297225 688884 848133 741618 091582 561157$
 $362135 750330 558036 085494 747230 138415 970602 017694 350758 458917 589235 971861 548843$
 $635060 827053 633582 882443 092203 262135 552296 661334 709021 156021 405492 515100 671199$
 $284761 072521 866782 927154 434480 887521.$

9. $\nu = 914 ?$

10. $\nu = 986 ?$

11. $\nu = 998$

$\Phi_2 = 23 \times 277 \times 4211 \times 1 385899 \times 240154 091459 652243 015812 929515 159070 212159 918817$
 $425875 611004 712759 052716 135663 441910 181493 025014 669780 274245 881010 561780 858639$
 $784499 969926 885693 756207 174479 909272 942309 784548 553831 369221 141895 942976 579419$
 $394048 307219 568666 715750 728448 387606 183250 921312 430705 694057 415487 884739 523892$
 $723969.$

Equations (5.4), (5.5) and (5.6) are a special case of equations

$$\Phi_1(\nu, \xi) = \Phi_1(\nu, 2^{2\xi+1}) = 2^{2\xi+1} + \Pi_1(\nu) = 2^{2\xi+1} + \frac{2^{\nu+3} + 1}{3}$$

$$\nu = 2\lambda, \lambda \in \mathbb{N}^* \quad (5.7)$$

$$\xi = 0, 1, 2, \dots, \frac{\nu - 2}{2} = \lambda - 1$$

$$\Phi_2(\nu, \xi) = \Phi_2(\nu, 2^{2\xi+1}) = -2^{2\xi+1} + \Pi_2(\nu) = -2^{2\xi+1} + \frac{5 \times 2^{\nu+1} - 1}{3}$$

$$\nu = 2\lambda, \lambda \in \mathbb{N}^* \quad (5.8)$$

$$\xi = 0, 1, 2, \dots, \frac{\nu - 2}{2} = \lambda - 1$$

$$\Phi_2(\nu, \xi) = 3 \times 2^{\nu+1} - \Phi_1(\nu, \xi) = \Phi_1^*(\nu, \xi) \quad (5.9)$$

for $\xi=0$. The general equations (5.7) and (5.8) or (5.9) give all possible variations of the method. For example, for $v=838$, a value of v that did not give a large prime number in the previous examples, from equation (5.7) for $\xi=1$, $2^{2\xi+1}=2^3$ we get

$$\begin{aligned}\Phi_1(838, 2^3) &= 3 \times 251 \times 124 \ 958179 \ 125661 \ 642577 \times 51 \ 945201 \ 394308 \ 356447 \ 274374 \ 943957 \\ &749268 \ 889249 \ 128703 \ 205379 \ 933327 \ 597692 \ 534177 \ 000888 \ 147927 \ 160249 \ 734500 \ 867000 \ 722765 \\ &431922 \ 957290 \ 626876 \ 299700 \ 840201 \ 468643 \ 187688 \ 745195 \ 339241 \ 792572 \ 155819 \ 582073 \ 320776 \\ &475981 \ 870379 \ 650986 \ 830637 \ 696975 \ 455178 \ 897139.\end{aligned}$$

For $v=914$ and $\xi=446$, $2^{2\xi+1}=2^{893}$ from equation (5.7) we get

$$\begin{aligned}\Phi_2(914, 2^{893}) &= \Phi_1^*(914, 2^{893}) = 664331 \times 2 \ 846040 \ 859139 \times 244156 \ 938515 \ 639118 \ 957128 \ 090989 \\ &142139 \ 910652 \ 427045 \ 074721 \ 151803 \ 096387 \ 973957 \ 589979 \ 330000 \ 475973 \ 653713 \ 544853 \ 032484 \\ &133776 \ 843427 \ 696009 \ 975250 \ 479123 \ 845904 \ 732248 \ 525911 \ 944098 \ 218552 \ 642872 \ 716830 \ 980122 \\ &426431 \ 711592 \ 329448 \ 776714 \ 147775 \ 168920 \ 208010 \ 693295 \ 505312 \ 315955 \ 821688 \ 402869.\end{aligned}$$

For $v=986$ and $\xi=2$, $2^{2\xi+1}=2^5$ from equation (5.7) we get

$$\begin{aligned}\Phi_1(986, 2^5) &= 29 \times 4253 \times 416 \ 339548 \ 376347 \times 33 \ 962793 \ 894011 \ 369474 \ 190517 \ 253893 \ 737579 \\ &763016 \ 839246 \ 389349 \ 317772 \ 361943 \ 015051 \ 440254 \ 338769 \ 687497 \ 378741 \ 306027 \ 784171 \ 795586 \\ &328198 \ 808732 \ 364434 \ 864222 \ 375712 \ 875623 \ 070283 \ 735838 \ 206731 \ 569184 \ 817760 \ 627665 \ 419282 \\ &986706 \ 989291 \ 001175 \ 496233 \ 325374 \ 140887 \ 038097 \ 481963 \ 826327 \ 954330 \ 255038 \ 364958 \ 385620 \\ &947977.\end{aligned}$$

For $v=66$ and $\xi=1, 2, 3, \dots, \frac{v-2}{2}=32$ from equation (5.7) we get

$$\Pi_1 = \Pi_1(66) = 196765270119568550571$$

$$\Pi_2 = \Pi_2(66) = 245956587649460688213 = 3^2 \times 27 \ 328509 \ 738828 \ 965357$$

$$\Phi_2(66, 2^1) = \Phi_1^*(66, 2^1) = 73 \times 3 \ 369268 \ 323965 \ 214907$$

$$\Phi_1(66, 2^3) = 1 \ 645337 \times 119 \ 589646 \ 449067$$

$$\Phi_2(66, 2^5) = 4140 \ 643009 \times 59400 \ 577909$$

$$\Phi_2(66, 2^7) = \Phi_1^*(66, 2^7) = 5 \times 13907 \times 3537 \ 162402 \ 379531$$

$$\Phi_2(66, 2^9) = \Phi_1^*(66, 2^9) = 13 \times 601 \times 25184 \ 342777 \ 367023$$

$$\Phi_2(66, 2^{11}) = \Phi_1^*(66, 2^{11}) = 5 \times 49 \ 191317 \ 529892 \ 137233$$

$$\Phi_1(66, 2^{13}) = 196\ 765270\ 119568\ 558763 \text{ is prime}$$

$$\Phi_2(66, 2^{15}) = \Phi_1^*(66, 2^{15}) = 5 \times 3259 \times 36269 \times 416167\ 836559$$

$$\Phi_2(66, 2^{17}) = \Phi_1^*(66, 2^{17}) = 157 \times 12\ 248491 \times 127901\ 671243$$

$$\Phi_1(66, 2^{19}) = 13 \times 19 \times 4643 \times 32083 \times 5347\ 833013$$

$$\Phi_1(66, 2^{21}) = 271 \times 5903 \times 123\ 000357\ 013771$$

$$\Phi_1(66, 2^{23}) = 7 \times 257879 \times 2\ 350441 \times 46\ 375123$$

$$\Phi_1(66, 2^{25}) = 634853 \times 309\ 938316\ 617551$$

$$\Phi_2(66, 2^{27}) = \Phi_1^*(66, 2^{27}) = 5 \times 49\ 191317\ 529865\ 294097$$

$$\Phi_1(66, 2^{29}) = 7^2 \times 107 \times 37529\ 137921\ 057681$$

$$\Phi_1(66, 2^{31}) = 13 \times 17\ 749367 \times 852750\ 974689$$

$$\Phi_1(66, 2^{33}) = 233 \times 1231 \times 5227 \times 131244\ 759703$$

$$\Phi_2(66, 2^{35}) = \Phi_1^*(66, 2^{35}) = 5 \times 14683 \times 3350\ 222537\ 834243$$

$$\Phi_2(66, 2^{37}) = \Phi_1^*(66, 2^{37}) = 73 \times 3\ 369268\ 322082\ 489517$$

$$\Phi_2(66, 2^{39}) = \Phi_1^*(66, 2^{39}) = 3^2 \times 27\ 328509\ 738828\ 965357$$

$$\Phi_1(66, 2^{41}) = 7 \times 37 \times 759711\ 476133\ 559097$$

$$\Phi_2(66, 2^{43}) = \Phi_1^*(66, 2^{43}) = 5 \times 41 \times 93\ 472763 \times 12835\ 698347$$

$$\Phi_1(66, 2^{45}) = 541 \times 13963 \times 26\ 047888\ 286741$$

$$\Phi_1(66, 2^{47}) = 7 \times 10061 \times 2793\ 891701\ 436337$$

$$\Phi_2(66, 2^{49}) = \Phi_1^*(66, 2^{49}) = 245\ 956024\ 699507\ 266901 \text{ is prime}$$

$$\Phi_2(66, 2^{51}) = \Phi_1^*(66, 2^{51}) = 5 \times 229 \times 8423 \times 25\ 502467\ 080379$$

$$\Phi_1(66, 2^{53}) = 7 \times 28\ 110611\ 045546\ 184509$$

$$\Phi_2(66, 2^{53}) = \Phi_1^*(66, 2^{53}) = 245\ 947580\ 450205\ 947221 \text{ is prime}$$

$$\Phi_1(66, 2^{55}) = 13 \times 19 \times 796766\ 392374\ 848237$$

$$\Phi_1(66, 2^{57}) = 601 \times 327636\ 248432\ 020643$$

$$\Phi_1(66, 2^{59}) = 7 \times 1447 \times 19482\ 844394\ 498171$$

$$\Phi_2(66, 2^{61}) = \Phi_1^*(66, 2^{61}) = 16183 \times 75731 \times 198808\ 531457$$

$$\Phi_2(66, 2^{63}) = \Phi_1^*(66, 2^{63}) = 5 \times 23 \times 41 \times 50208\ 529292\ 175167$$

$$\Phi_1(66, 2^{65}) = 7 \times 46507 \times 717\ 737600\ 997047.$$

Above, we saw the variation of the digits of the largest prime number for $\nu = 66$, as ξ takes values from $\xi = 0$ to $\xi = \frac{66-2}{2} = 32$. A similar variation arises for every $\nu \in \mathbb{N}^*$ and $\xi = 0, 1, 2, \dots, \frac{\nu-2}{2}$ in equations (5.7) and (5.9). From this distribution we come to an optimization of the method:

For a fixed computer network we choose a random time interval T ,

$$T = T(\nu) \quad (5.10)$$

and put for execution the following algorithmic procedure: The factorization of

$$\Phi_1(\nu, 2^{2\xi+1}) \text{ and } \Phi_2(\nu, 2^{2\xi+1}) = \Phi_1^*(\nu, 2^{2\xi+1}) = 3 \times 2^{\nu+1} - \Phi_1(\nu, 2^{2\xi+1})$$

is interrupted when the time interval $T = T(\nu)$ is exceeded. By this condition, the method gives only the largest primes for every $\nu \in \mathbb{N}^*$. Additionally, the required run time for the method is minimized for a fixed computer network and a fixed $\nu \in \mathbb{N}^*$. Next, we can see an example for a low computational power computer:

Example 5.1. For $\nu = 998$ and $T = 3s$ from equations (5.7) and (5.9) we get

$$\Pi_1 =$$

$$\begin{aligned} & 71433907145751154729895003270666787370760320780368907162916692558023403408329074832 \\ & 87989192104639054183964486117020978834580968571282093623989718383132383202623045183 \\ & 21615399028071640337491409458530278810203098332238796084493251170611036263071804194 \\ & 3047464318457694778440286554435082924558137112046251 \end{aligned}$$

$$\xi = 0, 1, 2, 3, \dots, \frac{998-2}{2} = 498.$$

1. $\xi = 0$, $\Phi_2(998, 2^1) = \Phi_1^*(998, 2^1) = 23 \times 277 \times 4211 \times 1385899 \times 240154 091459 652243 015812 929515 159070 212159 918817 425875 611004 712759 052716 135663 441910 181493 025014 669780 274245 881010 561780 858639 784499 969926 885693 756207 174479 909272 942309 784548 553831 369221 141895 942976 579419 394048 307219 568666 715750 728448 387606 183250 921312 430705 694057 415487 884739 523892 723969$ (288 digits)

T=0.6s.

2. $\xi = 1$, $\Phi_1(998, 2^3) = 97 \times 137 \times 9923 \times 1081681 \times 50080 644682 811942 296699 821816 279773 732717 457591 276109 457062 242788 762858 485087 453320 300320 950501 243057 574463 353237 805619 330431 715070 029378 288884 517452 798815 573077 642329 490846 332389 896689 408072 542886 142722 619257 370697 531361 739108 642652 205080 024148 915392 295551 861665 088028 145461 038179 893337$ (287 digits)

T=0.6s.

3. $\xi = 5$, $\Phi_2(998, 2^{11}) = \Phi_1^*(998, 2^{11}) = 3 \times 5 \times 23 \times 521 \times 150649 \times 1079767 946205 822485 067253 368196 496012 313340 598946 540309 265779 237716 776475 975549 878875 202589 444396 725173 878838 408179 822824 690856 508150 720159 542427 180957 684097 451829 502470 625266 660093 576980 953598 899215 688023 371196 514480 573388 038308 006579 994340 680318 041113 862815 165094 345737 124996 237057 863453$ (289 digits)

T=0.4s.

4. $\xi = 9$, $\Phi_1(998, 2^{19}) = 4536179 \times 356 426027 \times 4418 194305 492551 055438 288717 852823 131268 899490 581575 671658 140386 858327 881061 918039 334820 806178 879155 888881 005738 442642 836356 895420 852799 269710 636084 534157 398165 083264 381690 315538 696593 057882 325500 671380 059142 783057 514685 959697 790231 700820 951045 409420 029331 949677 361996 881725 184302 765883$ (286 digits)

T=0.6s.

5. $\xi = 13$, $\Phi_2(998, 2^{27}) = \Phi_1^*(998, 2^{27}) = 5 \times 1785847 678643 778868 247375 081766 669684 269008 019509 222679 072917 313950 585085 208226 870821 997298 026159 763545 991121 529255 244708 645242 142820 523405 997429 595783 095800 655761 295804 038497 570179 100843 728523 646325 697025 507745 830596 990211 233127 926527 590657 679510 485761 866079 614423 694610 071638 608770 731139 534251 168017$ (301 digits)

T=0.2s.

6. $\xi = 25$, $\Phi_2(998, 2^{51}) = \Phi_1^*(998, 2^{51}) = 5 \times 11661 167857 \times 16521 920251 \times 4880496 240331 \times 1899 231315 908657 097875 196652 086767 417005 191295 227630 772878 186939 449441 507048 325820 190068 477005 532913 880320 995340 694858 276546 924458 586740 967203 764636 699403 607786 693467 423245 127615 083671 437704 878429 984084 139551 880694 959912 111807 982429 926306 542433 872186 453701 258022 963289$ (268 digits)

T=2.6s.

7. $\xi = 29$, $\Phi_2(998, 2^{59}) = \Phi_1^*(998, 2^{59}) = 5^2 \times 11 \times 31 \times 41 \times 9 \ 612847 \times 8207 \ 166203 \times 323 \ 810490 \ 953628 \ 022064 \ 233800 \ 849387 \ 739635 \ 421729 \ 426241 \ 826230 \ 216969 \ 635899 \ 103869 \ 379857 \ 533713 \ 745016 \ 382039 \ 410786 \ 464059 \ 201318 \ 076449 \ 245745 \ 596630 \ 400732 \ 573528 \ 033619 \ 426479 \ 931615 \ 580202 \ 341503 \ 811059 \ 519577 \ 410940 \ 997169 \ 667449 \ 031997 \ 541911 \ 776463 \ 049074 \ 801526 \ 844609 \ 127839 \ 268529 \ 094893 \ 668906 \ 920738 \ 922813$ (279 digits)

T=1.1s.

8. $\xi = 38$, $\Phi_2(998, 2^{77}) = \Phi_1^*(998, 2^{77}) = 12689 \times 12791 \times 362 \ 367211 \times 63438 \ 358051 \times 73076 \ 239109 \times 32749 \ 562773 \ 009140 \ 773210 \ 602629 \ 062784 \ 951656 \ 363861 \ 516124 \ 158877 \ 033055 \ 872447 \ 866810 \ 557438 \ 212636 \ 674718 \ 860993 \ 747247 \ 501118 \ 202226 \ 318125 \ 517894 \ 724914 \ 745242 \ 508958 \ 914820 \ 486945 \ 219310 \ 864886 \ 943633 \ 570729 \ 820520 \ 875419 \ 898759 \ 256524 \ 768188 \ 352102 \ 694842 \ 651897 \ 783215 \ 819590 \ 508484 \ 531342 \ 501991$ (263 digits)

T=2.6s.

9. $\xi = 40$, $\Phi_2(998, 2^{81}) = \Phi_1^*(998, 2^{81}) = 2213 \times 46544 \ 703463 \times 86688 \ 749106 \ 696899 \ 807354 \ 911283 \ 542897 \ 445289 \ 843103 \ 778339 \ 244209 \ 182382 \ 057748 \ 453154 \ 185684 \ 585569 \ 053616 \ 848194 \ 373798 \ 960426 \ 348482 \ 124482 \ 392398 \ 508300 \ 413367 \ 195172 \ 964872 \ 513704 \ 529280 \ 109527 \ 456726 \ 374749 \ 787017 \ 361927 \ 178522 \ 932931 \ 830671 \ 329496 \ 439241 \ 623018 \ 611617 \ 086650 \ 436534 \ 193897 \ 918305 \ 285320 \ 911790 \ 616877 \ 929319$ (287 digits)

T=1.7s.

10. $\xi = 53$, $\Phi_1(998, 2^{107}) = 232109 \times 30 \ 776017 \ 795842 \ 106393 \ 933455 \ 088198 \ 556441 \ 482372 \ 842228 \ 826612 \ 891655 \ 454128 \ 622073 \ 391843 \ 846158 \ 439319 \ 908487 \ 194309 \ 417067 \ 485625 \ 196931 \ 532032 \ 671210 \ 912217 \ 922262 \ 312673 \ 714516 \ 124084 \ 732673 \ 674826 \ 399305 \ 513249 \ 870080 \ 803858 \ 480722 \ 552694 \ 304779 \ 051415 \ 207991 \ 530317 \ 060651 \ 323481 \ 394721 \ 722632 \ 287825 \ 024460 \ 688001 \ 173139 \ 722930 \ 671031$ (296 digits)

T=1.6s.

11. $\xi = 57$, $\Phi_2(998, 2^{115}) = \Phi_1^*(998, 2^{115}) = 5 \times 1 \ 785847 \ 678643 \ 778868 \ 247375 \ 081766 \ 669684 \ 269008 \ 019509 \ 222679 \ 072917 \ 313950 \ 585085 \ 208226 \ 870821 \ 997298 \ 026159 \ 763545 \ 991121 \ 529255 \ 244708 \ 645242 \ 142820 \ 523405 \ 997429 \ 595783 \ 095800 \ 655761 \ 295804 \ 038497 \ 570179 \ 100843 \ 728523 \ 646325 \ 697025 \ 507745 \ 830596 \ 990211 \ 233127 \ 926527 \ 590657 \ 679510 \ 485761 \ 866079 \ 614423 \ 686302 \ 396664 \ 953046 \ 525490 \ 740151 \ 259409$ (301 digits)

T=0.2s.

12. $\xi = 58$, $\Phi_1(998, 2^{117}) = 43691 \times 163 \ 497990 \ 766407 \ 623377 \ 572047 \ 494144 \ 760638 \ 942392 \ 667526 \ 280384 \ 785636 \ 762773 \ 576728 \ 225663 \ 941955 \ 761875 \ 549633 \ 887619 \ 063104 \ 919113 \ 291858 \ 299617 \ 112954 \ 467144 \ 027794 \ 812029 \ 553057 \ 221056 \ 584105 \ 513498 \ 906328 \ 611725 \ 849593 \ 401450$

827448 009505 902541 961615 264281 113052 176208 327617 656896 484824 281773 010317 594643
051590 431302 090753 (297 digits)

T=0.1s.

13. $\xi = 68$, $\Phi_2(998, 2^{137}) = \Phi_1^*(998, 2^{137}) = 105417 695783 \times 84 703410 816335 186156 806261$
492001 372948 895961 712695 154018 727889 782535 917962 496814 120959 304160 650293 671672
252582 708267 822424 427945 951346 549078 717762 234124 652165 456699 024462 611819 549559
013895 512611 534924 096905 108400 219574 841626 337387 088721 482649 840062 562405 309768
293988 524532 287345 514667 771427 (290 digits)

T=1.1s

14. $\xi = 73$, $\Phi_1(998, 2^{147}) = 97 \times 379 \times 5749 \times 12511 \times 2 701525 852731 749067 825970 181538$
841081 500468 570988 102390 004713 750684 212272 485899 120548 747955 954166 771149 039178
447733 176943 626139 060614 247893 193527 941427 464561 994282 082884 395184 241965 420437
569582 191909 943209 546029 373459 071393 266722 914761 826479 840486 134806 578495 819150
003168 800130 472275 972947 (289 digits)

T=0.1s

15. $\xi = 77$, $\Phi_2(998, 2^{155}) = \Phi_1^*(998, 2^{155}) = 5 \times 23 \times 593 \times 920 421803 063369 \times 142257 439151$
650269 759512 229987 535797 603859 632871 650964 499028 781520 394883 025379 626882 238970
651138 172615 583681 318636 124615 473290 390325 908383 034899 430684 465196 844839 590606
618826 905902 899124 401196 506691 143719 061249 533242 261455 674589 486954 446740 453073
599526 221890 802680 087230 458993 840559 (282 digits)

T=0.9s

16. $\xi = 83$, $\Phi_1(998, 2^{167}) = 114773 \times 155251 \times 400 894654 169098 163595 483363 914852 531928$
781693 381969 626300 885755 924123 202706 408560 360730 024076 936163 945689 446112 988849
825030 667217 976200 944962 868154 794364 056209 768391 777045 901886 394962 168952 671859
065166 558803 794909 960252 142707 475904 498255 382992 174430 898197 050208 510399 326290
198098 814515 285573 (291 digits)

T=0.4s

17. $\xi = 90$, $\Phi_1(998, 2^{181}) = 13 \times 239 \times 2029 \times 345814 471567 \times 3 276709 505948 663007 436071$
497798 167728 716237 262266 046137 399918 770194 945601 014426 083832 974746 796394 366422
444638 956063 948505 660872 339745 731543 151887 495733 528762 939605 565934 906885 716395
630970 274169 097750 084818 105248 219940 313650 115083 027840 487493 893745 793535 513039
225289 854656 163199 872403 (283 digits)

T=2.7s

18. $\xi = 95$, $\Phi_2(998, 2^{191}) = \Phi_1^*(998, 2^{191}) = 5 \times 17 \times 157 \times 2543 \times 263117 345854 197111 775236$
 642637 849621 102132 569635 056743 616085 430844 342072 473092 169620 260599 446899 578806
 313516 402000 281513 699879 810654 350772 703862 923881 305362 033902 791182 966354 140802
 049019 690624 171161 337402 154320 115975 545121 656939 078207 412594 447644 799532 466810
 505324 177709 194418 308158 570392 926619 (294 digits)

T=0.1s

19. $\xi = 96$, $\Phi_1(998, 2^{193}) = 13 \times 1129 \times 335029 \times 8 269211 \times 175 679317 658327 620139 490814$
 384990 040100 858556 390114 264495 835752 393980 708984 505460 993306 839388 894630 597492
 726644 756282 306765 232024 031693 331107 485972 274891 069941 204175 847675 178674 462934
 901401 328940 460822 002358 705536 882813 222690 476700 154781 597384 933080 860102 029870
 056922 793612 063768 927961 (285 digits)

T=0.3s

20/21. $\xi = 101$, $\Phi_1(998, 2^{203}) = 773 730593 \times 9232 400501 158825 799446 578593 625131 140544$
 202415 991183 014126 558203 447359 798054 291855 159964 176451 581841 722477 278492 490704
 493289 912566 869504 181528 102486 613687 773132 999515 783853 709246 722784 522123 832558
 993551 521661 213830 778384 182273 779765 239763 039295 444140 618834 911176 106897 416666
 287373 629641 164363 (292 digits)

T=1.5s

$\Phi_2(998, 2^{203}) = \Phi_1^*(998, 2^{203}) = 5 \times 1 785847 678643 778868 247375 081766 669684 269008 019509$
 222679 072917 313950 585085 208226 870821 997298 026159 763545 991121 529255 244708 645242
 142820 523405 997429 595783 095800 655761 295804 038497 570179 100843 728523 646325 697025
 507745 830596 990211 233127 926527 590655 108409 614947 481638 747284 346864 211474 573245
 941087 065741 529361 (301 digits)

T=0.2s

22/23. $\xi = 105$, $\Phi_1(998, 2^{211}) = 587 \times 720 547769 \times 16 888984 061692 358355 153236 651114$
 719240 611806 510989 941762 465501 533611 645216 696095 964090 405285 020017 953218 387778
 165387 576927 041108 914263 284391 156734 424355 035281 950141 293201 852464 222562 316879
 263965 425696 648530 916626 519821 543011 793530 618310 055052 637183 984261 934271 230666
 461625 182128 549715 877033 (290 digits)

T=0.5s

$\Phi_2(998, 2^{211}) = \Phi_1^*(998, 2^{211}) = 5 \times 37 \times 30491 \times 1 582963 939420 120308 648786 112132 928621$
 621628 730063 211101 789821 288825 665956 554505 557086 847335 568368 657783 813142 495087
 380422 087547 448933 113099 388148 736652 548603 757919 967348 839752 953400 605445 584318
 320182 824905 805386 818260 940278 551958 997672 853398 014378 684431 869765 404191 641547
 634034 247079 091754 668359 (295 digits)

T=0.2

24. $\xi = 123$, $\Phi_1(998, 2^{247}) = 1\ 371641 \times 16951\ 549039 \times 128216\ 263801 \times 2396\ 135860\ 767614$
 $483702\ 567537\ 826479\ 613108\ 014773\ 336354\ 737978\ 026710\ 481589\ 070522\ 440576\ 983324\ 884283$
 $105189\ 114843\ 934743\ 233073\ 871810\ 765202\ 096962\ 463805\ 602250\ 649971\ 484702\ 784738\ 167787$
 $009328\ 401038\ 173512\ 902558\ 630958\ 832427\ 970019\ 825426\ 400796\ 258991\ 194077\ 880602\ 439158$
 $915163\ 161887\ 574665\ 465621$ (274 digits)

T=2.0s

25. $\xi = 126$, $\Phi_2(998, 2^{253}) = \Phi_1^*(998, 2^{253}) = 107 \times 229 \times 557 \times 654244\ 322790\ 130219\ 004207$
 $626709\ 347972\ 072231\ 517142\ 195345\ 835320\ 100382\ 162959\ 860418\ 978785\ 508072\ 116827\ 875158$
 $506997\ 578377\ 075107\ 223247\ 035379\ 511073\ 680652\ 739397\ 497217\ 999306\ 022691\ 259692\ 441638$
 $920278\ 669033\ 252267\ 170826\ 591982\ 849056\ 108034\ 999793\ 424443\ 706688\ 901363\ 124456\ 987270$
 $278959\ 702729\ 922794\ 969298\ 190775\ 649151$ (294 digits)

T=0.5s

26. $\xi = 141$, $\Phi_1(998, 2^{283}) = 1879 \times 9151 \times 34511 \times 622103 \times 19\ 350368\ 061089\ 137024\ 255081$
 $876408\ 454341\ 134176\ 649502\ 367875\ 724581\ 146137\ 062600\ 954709\ 914300\ 522080\ 798917\ 742040$
 $679924\ 014966\ 680273\ 004127\ 182045\ 662741\ 209853\ 383601\ 157577\ 602353\ 579690\ 670975\ 286235$
 $843453\ 744059\ 341402\ 517087\ 478888\ 286731\ 306664\ 690147\ 302321\ 158859\ 165390\ 399832\ 045177$
 $744372\ 761706\ 710364\ 119787$ (284 digits)

T=0.6s

27. $\xi = 145$, $\Phi_2(998, 2^{291}) = \Phi_1^*(998, 2^{291}) = 5 \times 53 \times 1187 \times 28879 \times 42223 \times 24\ 611183 \times 3672$
 $479219 \times 257569\ 687091\ 602152\ 079137\ 833651\ 131746\ 881393\ 271929\ 762300\ 825308\ 494232$
 $053410\ 197389\ 850727\ 531290\ 943856\ 262690\ 752668\ 357924\ 658836\ 012815\ 730385\ 298925\ 272892$
 $377202\ 756649\ 576381\ 765451\ 465790\ 588151\ 386237\ 951574\ 352732\ 533550\ 835834\ 637309\ 905649$
 $834713\ 883058\ 580340\ 426909\ 925190\ 838256\ 544266\ 400827$ (270 digits)

T=1.0s

28/29. $\xi = 149$, $\Phi_1(998, 2^{299}) = 7\ 143390\ 714575\ 115472\ 989500\ 327066\ 678737\ 076032\ 078036$
 $890716\ 291669\ 255802\ 340340\ 832907\ 483287\ 989192\ 104639\ 054183\ 964486\ 117020\ 978834\ 580968$
 $571282\ 093623\ 989718\ 383132\ 383202\ 623045\ 183216\ 153990\ 280716\ 403374\ 914094\ 585302\ 788102$
 $030984\ 340905\ 949012\ 175554\ 840333\ 206835\ 407122\ 468781\ 661151\ 425820\ 096510\ 511231\ 625732$
 $806226\ 490203\ 744939$ (301 digits) is prime

T=0.3s

$\Phi_2(998, 2^{299}) = \Phi_1^*(998, 2^{299}) = 5^3 \times 11 \times 31 \times 41^2 \times 101 \times 251 \times 601 \times 1801 \times 3529 \times 4051 \times 8101 \times$
 $8431 \times 268501 \times 8598\ 989029 \times 2\ 014531\ 706153\ 918551\ 568531\ 289400\ 771225\ 903728\ 446888$
 $818738\ 054449\ 663123\ 056799\ 858197\ 716102\ 446457\ 294008\ 497198\ 921956\ 466935\ 821473\ 840872$

382211 084459 286634 758646 635538 904641 233614 993957 107131 230161 611503 584857 862662
299225 150156 806591 845847 074598 298104 560413 016403 (253 digits)

T=0.8s

30. $\xi = 176$, $\Phi_2(998, 2^{353}) = \Phi_1^*(998, 2^{353}) = 23 \times 388227 756226 908449 618994 582992 754279$
188914 786849 831017 189764 633467 518496 784397 145830 868977 831773 861640 432852 506359
835806 227226 552787 070305 651615 129518 064304 490382 890392 182282 080473 717574 723592
096229 586791 287745 959134 712343 206475 830707 192755 437013 192996 026942 738118 933232
281727 325378 488296 376778 157427 (300 digits)

T=0.5s

31. $\xi = 185$, $\Phi_1(998, 2^{371}) = 2099 \times 12 937471 \times 680 377837 \times 386627 212955 742182 618863 921271$
402758 475803 518040 436431 392999 790429 241969 001800 633833 709644 148165 858361 909113
085591 159491 017541 821555 845870 762554 472234 194100 606449 540385 960430 250521 051832
567287 314534 455016 244516 023635 797172 266397 392463 416148 563961 924352 900518 142198
481657 007282 974163 (282 digits)

T=1.7s

32. $\xi = 196$, $\Phi_2(998, 2^{393}) = \Phi_1^*(998, 2^{393}) = 5 402543 \times 1 652784 326421 630395 396552 217878$
385867 793192 964414 371786 650210 200965 161300 158302 183640 183241 508822 570728 258082
841039 159437 181011 000579 285908 504041 148569 382789 415800 388635 535614 959639 466829
753219 118784 455475 999024 511133 864710 625928 040615 180536 068168 167287 105977 962653
037457 737160 438527 849330 239547 (295 digits)

T=0.5s

33. $\xi = 197$, $\Phi_1(998, 2^{395}) = 11 \times 1999 \times 27091 \times 282 372962 966749 \times 42 466921 406365 948875$
948770 737359 680169 193512 590162 300543 819779 157936 035393 558127 029345 608326 557832
279148 245930 623123 641140 114550 223537 988714 596718 353365 008099 378198 523221 591670
666786 220417 262224 405322 567840 139735 569537 878617 569393 087598 107572 123285 078427
876051 267706 022614 262169 (278 digits)

T=2.6s

34/35. $\xi = 199$, $\Phi_1(998, 2^{399}) = 201575 740129 \times 2 480055 715723 \times 14 289094 363187 213691$
208680 611171 705945 303353 242452 569850 169220 526954 050535 261843 940156 925970 575589
033037 189974 594162 587028 204034 474059 758365 011979 903011 741983 421318 377292 028446
684764 288036 524278 562883 228210 081177 068286 210999 538629 937470 406625 038023 246173
066205 804579 013580 057217 (278 digits)

T=2.5s

$$\Phi_2(998, 2^{399}) = \Phi_1^*(998, 2^{399}) = 5^3 \times 11 \times 17 \times 31 \times 41 \times 101 \times 157 \times 251 \times 401 \times 601 \times 1801 \times 4051 \times 8101 \times 15671 \times 61681 \times 268501 \times 340801 \times 2787601 \times 18004897 \times 3173389601 \times 793803149263 \times 474063717704894849562835374581495827311871974257278381183325961420859899717498041765514624204329617053587500499410353583821054381178850330054652795756561865556583500734550291163982011165238195274533641100352256127741 (216 digits)$$

T=1.6s

$$36. \xi = 201, \Phi_1(998, 2^{403}) = 262957 \times 3887239 \times 28552837 \times 244753639896519737750246725613155331950850159334731536596181252980547298934418475724098648820062993092677736512139629575073122323492039685203991554440481720353819310553897215152630195141276654005494101148387952195777353069857283492453953255273438038066180257642229578928507868082909 (282 digits)$$

T=0.9s

$$37. \xi = 205, \Phi_2(998, 2^{411}) = \Phi_1^*(998, 2^{411}) = 5 \times 1785847678643778868247375081766669684269008019509222679072917313950585085208226870821997298026159763545991121529255244708645242142820523405997429595783095800655761295804038497569121411293664125888002632532655312166871542718490277975672152585240461219783767926294120957906527182852624791899496904724753 (301 digits)$$

T=0.2s

$$38. \xi = 214, \Phi_2(998, 2^{429}) = \Phi_1^*(998, 2^{429}) = 11 \times 31 \times 3539 \times 45953 \times 161014729636910389111812999030059341331702777776625204197132458005105750940625571667284320311378697556043726460762240869632050174893072751159659708275344562220222729353141825480373379280680480875360857256748360641295229595555109930503027083453002486844005086833270345095169720212907905003483 (291 digits)$$

T=0.2s

$$39. \xi = 225, \Phi_2(998, 2^{451}) = \Phi_1^*(998, 2^{451}) = 5 \times 83 \times 157 \times 593 \times 1547469837457 \times 149344688969630053142970351119285083892632041523756578296865620098073334610600116811906551196474836660748584447955749792450847018523472038962609946451521773976632269231508876581693377964213322116002873927538556522954025942807762980646429619608682337730938114538150182192249431760423 (282 digits)$$

T=1.6s

$$40. \xi = 235, \Phi_2(998, 2^{471}) = \Phi_1^*(998, 2^{471}) = 5 \times 17 \times 105049863449634051073375004809804099074647530559366039945465724350034416776954521813058664589774103737999477737015014394626190714283560200352789976222535047097325991928033606653223274377990621650478979463295837733811331295522321593591242336018824465544275854610595076449457186855308021508224252444929 (300 digits)$$

T=0.2s

41. $\xi = 239$, $\Phi_1(998, 2^{479}) = 137 \times 2 \ 434997 \times 827 \ 931791 \times 25 \ 863711 \ 562732 \ 751591 \ 526226$
 $429290 \ 028370 \ 217018 \ 792199 \ 675768 \ 141948 \ 292179 \ 277815 \ 856585 \ 213443 \ 474425 \ 000554 \ 419832$
 $365597 \ 844897 \ 342275 \ 207952 \ 315785 \ 919467 \ 304239 \ 926018 \ 076282 \ 134454 \ 047167 \ 786222 \ 235401$
 $711152 \ 690088 \ 464424 \ 367353 \ 699351 \ 674265 \ 583706 \ 525493 \ 638382 \ 489694 \ 557908 \ 274402 \ 227935$
 $227883 \ 791745 \ 829585 \ 571561$ (284 digits)

T=1.2s

42. $\xi = 240$, $\Phi_1(998, 2^{481}) = 13 \times 451313 \times 32 \ 388773 \times 37591 \ 418449 \ 933545 \ 132030 \ 943408$
 $610794 \ 177762 \ 760160 \ 064590 \ 467852 \ 865002 \ 096742 \ 764103 \ 807119 \ 613757 \ 199281 \ 228573 \ 494102$
 $705882 \ 841103 \ 511342 \ 128120 \ 867188 \ 414940 \ 162429 \ 183971 \ 386811 \ 945315 \ 764966 \ 480741 \ 277831$
 $641063 \ 626381 \ 603864 \ 827290 \ 129144 \ 053231 \ 083021 \ 001899 \ 477302 \ 962558 \ 981548 \ 628055 \ 404202$
 $683736 \ 887772 \ 673100 \ 322619$ (287 digits)

T=0.5s

43/44. $\xi = 244$, $\Phi_1(998, 2^{489}) = 109 \times 15737 \times 1 \ 810423 \times 63 \ 897091 \times 35999 \ 363077 \ 864234 \ 579333$
 $215835 \ 234616 \ 573581 \ 578631 \ 320360 \ 833997 \ 586420 \ 240784 \ 265629 \ 453544 \ 798943 \ 077771 \ 079413$
 $782788 \ 321913 \ 294236 \ 465957 \ 355591 \ 614068 \ 051739 \ 603078 \ 286428 \ 192972 \ 852372 \ 442468 \ 665132$
 $012334 \ 653651 \ 970015 \ 916188 \ 872946 \ 706635 \ 088401 \ 901774 \ 944026 \ 904859 \ 962958 \ 461711 \ 445028$
 $659621 \ 722389 \ 359879 \ 084227$ (281 digits)

T=0.4s

$\Phi_2(998, 2^{489}) = \Phi_1^*(998, 2^{489}) = 11 \times 29 \times 31 \times 6131 \times 15859 \ 071337 \times 46487 \ 196589 \times 199 \ 765099$
 $267706 \ 307384 \ 558500 \ 862817 \ 361682 \ 347934 \ 156991 \ 993085 \ 611675 \ 011340 \ 665792 \ 527850 \ 931325$
 $277078 \ 135380 \ 228309 \ 779364 \ 438194 \ 606702 \ 523877 \ 725780 \ 747131 \ 984481 \ 430096 \ 817822 \ 024230$
 $157706 \ 713138 \ 235841 \ 861068 \ 101965 \ 998665 \ 309826 \ 441512 \ 161687 \ 509188 \ 054991 \ 629800 \ 107672$
 $152543 \ 620661 \ 007412 \ 372772 \ 259323$ (273 digits)

T=1.7s

45. $\xi = 250$, $\Phi_1(998, 2^{501}) = 701 \times 88993 \times 114506 \ 605306 \ 822613 \ 786971 \ 308969 \ 205959 \ 876278$
 $590409 \ 848400 \ 237087 \ 189130 \ 125532 \ 815439 \ 145413 \ 046065 \ 959549 \ 000002 \ 035206 \ 091992 \ 408687$
 $172561 \ 258064 \ 328892 \ 687270 \ 749983 \ 315938 \ 949995 \ 433402 \ 489002 \ 599979 \ 639067 \ 275489 \ 260978$
 $578588 \ 631521 \ 276723 \ 222191 \ 236183 \ 570319 \ 390598 \ 743975 \ 491594 \ 455620 \ 068289 \ 121723 \ 044383$
 $566285 \ 939440 \ 760871$ (294 digits)

T=0.2s

46. $\xi = 278$, $\Phi_2(998, 2^{557}) = \Phi_1^*(998, 2^{557}) = 229 \times 727 \times 53 \ 634535 \ 617563 \ 921488 \ 902022 \ 481775$
 $006585 \ 327271 \ 238181 \ 156006 \ 106248 \ 504369 \ 367599 \ 341280 \ 215457 \ 352943 \ 728782 \ 024170 \ 369392$
 $948686 \ 794106 \ 462586 \ 054444 \ 869954 \ 794455 \ 207116 \ 935716 \ 897118 \ 296816 \ 058532 \ 693088 \ 478493$

873803 596149 649557 250353 168036 287281 695021 634012 300003 716337 092657 262753 514640
377523 759265 284891 362434 694327 (296 digits)

T=0.2s

47. $\xi = 279$, $\Phi_2(998, 2^{559}) = \Phi_1^*(998, 2^{559}) = 5^2 \times 11 \times 17 \times 31 \times 41 \times 827 \times 61681 \times 308 397511 \times 95$
525533 582727 010358 730178 212996 182198 046190 995203 859912 007761 218911 979801 056527
425688 304386 103373 878566 313258 491033 358792 611575 525127 777226 451065 150174 051526
998643 029803 956912 334773 819612 574956 958053 168970 683768 780685 731223 766665 224274
355223 675088 259411 620017 146224 397018 314357 (278 digits)

T=1.5s

48. $\xi = 290$, $\Phi_2(998, 2^{581}) = \Phi_1^*(998, 2^{581}) = 4271 \times 812 807357 \times 47 351674 899017 \times 54320 264421$
708924 159885 550778 684086 375546 537121 076070 874827 461311 350281 259344 523914 762800
663168 792865 931037 276353 338365 029872 927679 239415 307317 576197 627916 475652 149701
055322 844380 011460 685315 101399 407826 653375 607407 270811 361373 662256 367298 430634
913311 553196 800451 968236 359839 (275 digits)

T=1.2s

49. $\xi = 291$, $\Phi_1(998, 2^{583}) = 349 \times 2907 294413 \times 23810 368063 \times 73546 783547 \times 4020 315822$
611107 366990 471894 187419 248314 220588 895788 829779 811649 899365 757256 877628 752645
126017 473183 589056 566855 137954 948482 141605 536988 074000 928848 554254 420994 950268
969954 835038 601519 213871 482150 608991 854938 042192 490200 506235 111973 470762 341180
295165 477898 598875 810087 (268 digits)

T=2.9s

50. $\xi = 295$, $\Phi_2(998, 2^{591}) = \Phi_1^*(998, 2^{591}) = 5 \times 17 \times 1406 097989 \times 74 710201 046759 374231$
190230 946133 654611 639964 844132 950356 894169 734876 433833 627025 984643 993818 964279
692174 054974 069815 268490 576987 062960 378081 613146 189500 797838 615894 934933 459515
736005 792665 911397 038134 695613 281592 216260 281415 828570 054065 928088 832312 687295
573298 633464 464089 053383 199273 268621 (290 digits)

T=0.7s

51. $\xi = 310$, $\Phi_1(998, 2^{621}) = 36855 430043 \times 193 821933 599493 272177 567582 943172 082662$
436241 377515 669660 033690 026704 115775 738622 638000 110666 749358 113953 414387 433082
786982 344762 852751 929213 620383 019188 075853 104891 533278 633490 641964 432434 239056
108392 905102 005458 066852 263624 668794 880022 911325 282562 473719 604069 735818 320825
292180 309996 201521 (291 digits)

T=0.8s

52. $\xi = 322$, $\Phi_2(998, 2^{645}) = \Phi_1^*(998, 2^{645}) = 277 \times 32235\ 517665\ 050160\ 076667\ 420248\ 495842\ 676335\ 884828\ 686329\ 947164\ 572453\ 981680\ 238415\ 647487\ 761684\ 079894\ 580569\ 422221\ 883206\ 413211\ 671405\ 105402\ 928644\ 839944\ 950886\ 654321\ 477837\ 646010\ 323536\ 862980\ 306233\ 004600\ 826793\ 019239\ 143486\ 688083\ 316231\ 173760\ 641133\ 887614\ 708958\ 874886\ 812632\ 805841\ 496062\ 168876\ 045362\ 639002\ 611983\ 749953$ (299 digits)

T=0.1s

53. $\xi = 324$, $\Phi_1(998, 2^{649}) = 13 \times 572821 \times 959272\ 780552\ 485045\ 736465\ 174053\ 792712\ 138163\ 187511\ 643215\ 203845\ 966621\ 918317\ 191168\ 120755\ 133089\ 918764\ 937601\ 525785\ 016559\ 368378\ 137849\ 325617\ 522371\ 822700\ 381247\ 853164\ 132156\ 968593\ 265532\ 025688\ 149083\ 061002\ 134066\ 754678\ 057925\ 758911\ 347548\ 291107\ 226398\ 649014\ 798037\ 634476\ 921719\ 554770\ 271195\ 412905\ 882872\ 170718\ 999931$ (294 digits)

T=0.5s

54. $\xi = 328$, $\Phi_2(998, 2^{657}) = \Phi_1^*(998, 2^{657}) = 29 \times 661 \times 465\ 816599\ 364541\ 412762\ 109416\ 705793\ 125428\ 819453\ 155934\ 759004\ 882183\ 199589\ 202673\ 125064\ 119671\ 682932\ 380344\ 187456\ 706110\ 367950\ 710601\ 934115\ 989371\ 821453\ 948705\ 306139\ 900675\ 983648\ 459184\ 031151\ 052206\ 260865\ 961407\ 827115\ 708474\ 752170\ 303864\ 828688\ 979588\ 262185\ 397208\ 543982\ 696577\ 825655\ 462101\ 679409\ 961161\ 593780\ 039205\ 203189$ (297 digits)

T=0.1s

55. $\xi = 335$, $\Phi_2(998, 2^{671}) = \Phi_1^*(998, 2^{671}) = 5 \times 17 \times 1579 \times 66\ 529362\ 539350\ 254004\ 670680\ 690186\ 256538\ 725478\ 504981\ 659243\ 486842\ 526937\ 566039\ 869868\ 152665\ 398726\ 899368\ 982097\ 299007\ 797871\ 695768\ 173511\ 983186\ 737589\ 118822\ 301408\ 609592\ 427860\ 122784\ 039449\ 766451\ 405543\ 530905\ 392162\ 706272\ 467897\ 814273\ 122233\ 108312\ 621456\ 492962\ 446789\ 429307\ 147288\ 453987\ 268908\ 221454\ 645762\ 134691\ 941251$ (296 digits)

T=0.1s

56. $\xi = 336$, $\Phi_1(998, 2^{673}) = 13 \times 643 \times 947 \times 4817 \times 15\ 563039 \times 6452\ 647781 \times 1865\ 482111\ 694196\ 984405\ 721028\ 109267\ 512169\ 662287\ 520722\ 309037\ 678721\ 775415\ 465389\ 736977\ 685952\ 209537\ 805133\ 533760\ 037167\ 161890\ 202547\ 864033\ 575096\ 403601\ 034516\ 354194\ 440294\ 113199\ 035912\ 051501\ 044428\ 259660\ 643442\ 949856\ 321344\ 053411\ 279634\ 717439\ 526154\ 204509\ 705020\ 485323\ 674761\ 245712\ 042546\ 515597$ (274 digits)

T=0.7s

57. $\xi = 337$, $\Phi_2(998, 2^{675}) = \Phi_1^*(998, 2^{675}) = 5 \times 1\ 785847\ 678643\ 778868\ 247375\ 081766\ 669684\ 269008\ 019509\ 222679\ 072917\ 313950\ 585085\ 208226\ 870821\ 997298\ 026159\ 732193\ 137933\ 322263\ 648242\ 402673\ 172327\ 054463\ 714020\ 815840\ 943906\ 796237\ 351353\ 694071\ 648607\ 168608\ 152146\ 800183\ 721008\ 933505\ 037501\ 681553\ 156304\ 073250\ 246384\ 262689\ 411986\ 276880\ 407170\ 664997\ 779636\ 981338\ 120611\ 558293\ 836049$ (301 digits)

T=0.3s

58. $\xi = 338$, $\Phi_2(998, 2^{677}) = \Phi_1^*(998, 2^{677}) = 107 \times 83450 826104 849479 824643 695409 657461 881729 346706 038442 947332 584764 046032 019076 022001 027911 122717 740099 746652 743050 881254 377334 872002 179796 115130 582991 331225 477461 005514 143027 751583 708967 356215 713166 252007 439756 198982 044653 220833 295019 542689 907113 373395 304172 092776 241876 677739 817712 163038 674315 014463 (299 digits)$

T=0.2s

59. $\xi = 352$, $\Phi_2(998, 2^{705}) = \Phi_1^*(998, 2^{705}) = 23 \times 10651 \times 13679 \times 37 146433 \times 71013 472483 \times 877312 715381 \times 1151 408834 272769 616264 028242 815038 516160 668354 436413 248269 807293 123140 307312 762024 536657 068208 059014 781694 377143 542542 796647 355483 353174 367781 256346 523370 543923 764970 333121 124603 680203 764935 092487 648250 393836 749893 110771 633048 079830 195769 053937 473038 020613 683577 (262 digits)$

T=2.9s

60. $\xi = 359$, $\Phi_2(998, 2^{719}) = \Phi_1^*(998, 2^{719}) = 5^2 \times 11 \times 17 \times 31 \times 41 \times 30497 \times 61681 \times 14 160555 016603 \times 56 415501 183158 349160 686674 861979 191473 446248 684175 761835 579034 727284 183059 885039 649800 383964 528477 523288 742866 508024 938336 809504 120750 319466 108619 183801 412427 589818 862624 067559 950557 831114 436890 657855 493642 551670 132567 915170 686046 491491 868627 512790 813166 283642 607570 231603 (272 digits)$

T=2.6s

61. $\xi = 361$, $\Phi_1(998, 2^{723}) = 97 \times 149 \times 494 249686 194915 621185 186489 107221 942646 926733 414300 886756 498253 359326 115051 055662 030866 561282 049052 277990 335001 485386 369573 505392 926243 163853 032142 313843 098458 228170 223858 991616 355435 276080 622915 380214 117152 667048 483900 878820 826383 015214 525959 955087 003729 585719 578697 235091 619395 197652 142456 697503 (297 digits)$

T=0.1s

62/63. $\xi = 366$, $\Phi_1(998, 2^{733}) = 13 \times 197 \times 347 \times 8 038321 119806 536613 815411 540055 699983 318872 061229 786541 293498 302291 342359 773579 449301 270927 620941 258502 572911 327940 378630 289589 009744 357063 355128 733522 983279 225000 677642 907764 056353 997324 939563 558957 447586 982342 515974 063254 044459 282552 480876 926346 614092 110408 702983 976784 026113 890618 854881 100529 (295 digits)$

T=0.2s

$\Phi_2(998, 2^{733}) = \Phi_1^*(998, 2^{733}) = 28949 \times 308 447213 831873 098940 788124 247239 919214 654741 011645 079117 226383 286224 927494 598580 570166 953981 115716 140254 037527 107340 838079 506065 455187 470916 309861 335942 285540 184738 161506 669113 646204 354305 663259 099907$

946899 335904 477867 595817 341664 358017 561769 321520 790058 421222 113323 078307 182041
 732080 360122 475329 (297 digits)

T=0.2s

64. $\xi = 398$, $\Phi_2(998, 2^{797}) = \Phi_1^*(998, 2^{797}) = 29 \times 163 \times 2463029 \times 766936293641854859996$
 $241608214948660659056423226601016664353466779874594217421947834134116883884398$
 $435553582671602574262856904998834876720809438464792987481224305027834865905116$
 $302509209202130027079252960983420913328570423486026674436711658835296797774424$
 $285951284461891051756446211637340327$ (291 digits)

T=1.1s

65. $\xi = 399$, $\Phi_2(998, 2^{799}) = \Phi_1^*(998, 2^{799}) = 5^3 \times 11 \times 17 \times 31 \times 41 \times 101 \times 251 \times 401 \times 601 \times 727 \times$
 $1381 \times 1801 \times 4051 \times 8101 \times 18341 \times 61681 \times 268501 \times 340801 \times 2787601 \times 3096427 \times 22236047 \times$
 $385568933 \times 3173389601 \times 8409382921 \times 4055093275862234632967152290766093157926308$
 $080381263801953434067536396366528848607438580599723331031593797701209811504204$
 $951986391814164791656984656394323556901357576913272343943139972622741238478271$
 (199 digits)

T=0.7s

66. $\xi = 400$, $\Phi_1(998, 2^{801}) = 10125793 \times 705464817874028777103136547139239241516791$
 $038295656519572509786821753723440865880133197531750869620157217229103275345580$
 $361409387467996717702283786320155499051237867793400561986760758488567002437780$
 $775768785283778019178464352524265035566971129062738206965345676132876866458749$
 051176325397992971 (294 digits)

T=0.7s

67. $\xi = 401$, $\Phi_1(998, 2^{803}) = 61927 \times 115351796705396926590816611931252583478547839$
 $844282634655186309685885048196539475868398516998514641630476900437326273659261$
 $308974585958219848561852515693568830219268210972620282014040948021816808596342$
 $386203911540406554848540517445097906684596465781569352560895680105144058973959$
 975176195334287517 (297 digits)

T=0.1s

68. $\xi = 408$, $\Phi_2(998, 2^{817}) = \Phi_1^*(998, 2^{817}) = 142699 \times 1846951 \times 33879587548349245200557671$
 $253217043154307156362321278579421774944084759667412391141840582765349102067532$
 $624290317216351463868244562251789887979126573411373583249089114801281543585495$
 $422391457605383196910025888366843381947903287967919488663050817933454479804737$
 $156647874763758725791514916009$ (290 digits)

T=1.0s

69. $\xi = 415$, $\Phi_2(998, 2^{831}) = \Phi_1^*(998, 2^{831}) = 5 \times 17 \times 13 \times 738939 \times 324012083 \times 1214035279 \times 19437919695183749611689831947804714931428224339795663392913311810739121573995786880746093626329331972413585305942189476499679941948766165414860971838232913112919122610133519492893365897765604686820606017425865907912209695993061042084194243462261483785539732509260071971982063 (275 digits)$

T=2.6s

70. $\xi = 422$, $\Phi_2(998, 2^{845}) = \Phi_1^*(998, 2^{845}) = 30047 \times 1307821 \times 227229647255682705642525444617843660928229364164101761305625487466711149585594347205782604644883381310419363594017430217640792216224451435993727300680274695078127345861091023469036113840120003285534680970159441873792882422851664415544249782082010908086736203950105875888687234292184995935063 (291 digits)$

T=0.4s

71. $\xi = 459$, $\Phi_1(998, 2^{919}) = 71 \times 100611136825001626380133869841160689830515534524580692750171927220523525437326718001073473506106358207698232001259282480038994356902384918601202973028509572215872738695569479384835982727925613832554139247839295999079475912287843371686635199635230461199158862424572210574543796146342802140426645878909 (300 digits)$

T=0.2s

72. $\xi = 467$, $\Phi_2(998, 2^{935}) = \Phi_1^*(998, 2^{935}) = 5 \times 17 \times 193 \times 4951 \times 2994643 \times 167360569151 \times 316897488881 \times 692194682407018939975765511703243084765560211784736983754900808956805839935450928887642673596592964383170982790765412920328379441353607935129845750509440143843391470934822922000125919848831773693595718121554122562347408091408627813895611285464692861131672925217043 (264 digits)$

T=2.2s

73. $\xi = 468$, $\Phi_1(998, 2^{937}) = 13 \times 549491593428855036473171714370380519606693723439730809914939714488364789382325326798351371577784997770471325524395625210220093964028488755358258700478026214371894736515569214744127874805730163161771272752698160788635564042299582463334676254458952750916938297354237223522572830236240988957823368523671 (300 digits)$

T=0.2s

74. $\xi = 470$, $\Phi_2(998, 2^{941}) = \Phi_1^*(998, 2^{941}) = 773 \times 457035727 \times 25274628064090325719460651342540761864690763486298197189937998391605967426911349278035633411559301717752687815184121224633404134985360068405148401351842766264662482564264083436142575797378518119852203977674522509813997648730863599746227144490943614288832177284925162410943526337233834536991 (290 digits)$

T=0.9

75. $\xi = 475$, $\Phi_2(998, 2^{951}) = \Phi_1^*(998, 2^{951}) = 5 \times 17 \times 105049 863449 633827 146122 904626 001701 248765 456221 399337 826282 415641 221758 966603 942471 406762 787703 734583 241530 080529 278519 053667 014635 853778 956423 588283 793168 007616 098451 399793 715578 091490 967719 433477 442424 900079 883426 811023 032581 199593 582414 723910 870799 328163 944101 834339 682451 529432 582488 249488 965889$ (300 digits)

T=0.1s

76. $\xi = 477$, $\Phi_1(998, 2^{955}) = 11 \times 137 \times 1151 \times 811 560341 \times 545823 333697 \times 9297 001807 780365 378157 339991 931775 374994 174254 862155 487110 020205 625810 301420 380556 294833 009868 905666 872651 884681 274096 795146 081058 186377 800015 494430 395149 833924 359413 745004 673172 691476 704437 295834 061266 627712 469047 819535 451890 375103 771868 768002 035400 864495 188032 572747 188171$ (274 digits)

T=1.6s

77. $\xi = 480$, $\Phi_1(998, 2^{961}) = 13 \times 705161 \times 779242 745175 008707 887255 024600 207012 392236 279773 603602 582286 366120 216752 873275 462643 339829 137717 953152 185595 437449 100676 413317 276005 491813 350221 773819 196464 704314 714605 297071 348641 204952 830267 337342 699086 693740 590320 878996 501931 735825 212243 467971 217296 301264 717544 619220 781889 038242 175762 602271$ (294 digits)

T=0.5s

78. $\xi = 483$, $\Phi_2(998, 2^{967}) = \Phi_1^*(998, 2^{967}) = 5 \times 17 \times 37 \times 53 \times 53 569537 702681 669902 976725 719459 552303 428718 803455 358056 168375 116620 172178 167663 087211 710405 946546 415676 625338 005442 566911 275166 851365 409528 398787 779065 282595 019268 928339 016361 739010 447221 894845 540400 264801 154160 168251 483698 796430 981716 466960 686431 808676 862342 923285 876071 696115 531569 104470 746521$ (296 digits)

T=0.2s

By increasing the values of $v \in \mathbb{N}^*$ the method gives even bigger primes.

As ξ takes values from $\xi = 0$ to $\xi = \frac{v-2}{2}$ we have neighborhoods of \mathbb{N} where the method gives big primes. We have such a neighborhood for small values of $\xi \in \mathbb{N}$. Next, we have an example for ten consecutive even numbers.

Example 5.2.

1.v=900

$\Pi_1 =$

22540566661788383844366497489771375211470819244207851610792990251804753175753594004

24180291787337335017277418637897082363231349587211458292268966308108542194809116518
 39575502578658780875275281515957006040836773719103158738240666341904382676807268896
 63553646118303487929003

$$\xi_{\max} = 449$$

$\xi = 0, \Phi_2(900, 2^1) = \Phi_1^*(900, 2^1) = 17 \times 1303 \times 14 \ 612289 \ 234017 \times 87 \ 048891 \ 637396 \ 895102$
 119311 765818 573790 576976 052955 671628 663425 575545 670256 411482 623556 091041 420326
 465040 293560 953469 136495 015271 490122 045045 761776 713210 784123 751846 457476 167032
 775232 117142 820499 562144 984787 267535 228165 750170 673557 613732 612894 777923 633653
 (254 digits)

T=2.9s

2.v=902

$$\Pi_1 =$$

90162266647153535377465989959085500845883276976831406443171961007219012703014376016
 96721167149349340069109674551588329452925398348845833169075865232434168779236466073
 58302010314635123501101126063828024163347094876412634952962665367617530707229075586
 54214584473213951716011

$$\xi_{\max} = 450$$

$\xi = 13, \Phi_1(902, 2^{27}) = 79 \times 229 \times 1 \ 460479 \times 3412 \ 454982 \ 716461 \ 315441 \ 492312 \ 197972 \ 257894$
 940472 842179 546378 500415 486322 912154 774592 223763 649210 130907 126821 500568 874032
 042527 791656 219277 508682 946942 316385 551856 434126 811018 680404 462255 372535 998280
 156025 411962 803812 239151 596564 022946 652154 073044 186875 329433 668351 (262 digits)

T=0.3s

3.v=904

$$\Pi_1 =$$

36064906658861414150986395983634200338353310790732562577268784402887605081205750406
 78688466859739736027643869820635331781170159339538333267630346092973667511694586429
 43320804125854049400440450425531209665338837950565053981185066147047012282891630234
 616858337892855806864043

$$\xi_{\max} = 451$$

$\xi = 0, \Phi_1(904, 2^1) = 3^3 \times 5 \times 55349 \ 376221 \times 269 \ 949933 \ 997411 \times 178794 \ 883900 \ 768710 \ 779431$
 681652 203639 002785 735652 238602 781429 372818 766589 763466 986895 723986 365238 027066
 477543 364144 304665 741374 114296 962532 152246 702412 638211 290697 066742 179336 399463
 717647 854627 879305 921306 176521 561550 312327 666557 219221 818976 363757 (246 digits)

T=3.8s

4.v=906

$\Pi_1 =$

14425962663544565660394558393453680135341324316293025030907513761155042032482300162
 71475386743895894411057547928254132712468063735815333307052138437189467004677834571
 77328321650341619760176180170212483866135535180226021592474026458818804913156652093
 8467433351571423227456171

$\zeta_{\max} = 452$

$$\xi = 0, \Phi_1(906, 2^1) = 2137 \times 7904 \ 941489 \times 85 \ 396804 \ 898645 \ 518273 \ 535449 \ 436485 \ 582682 \ 070411 \\ 560076 \ 424907 \ 006632 \ 243214 \ 309176 \ 544188 \ 795999 \ 633591 \ 556075 \ 663051 \ 094702 \ 367573 \ 582291 \\ 715723 \ 564664 \ 172471 \ 347636 \ 830470 \ 189043 \ 981602 \ 095079 \ 931244 \ 982740 \ 899398 \ 967082 \ 280998 \\ 040450 \ 772949 \ 380551 \ 511813 \ 090613 \ 292727 \ 464205 \ 777687 \ 764915 \ 232261 \ (260 \text{ digits})$$

T=0.4s

5.v=908

$\Pi_1 =$

57703850654178262641578233573814720541365297265172100123630055044620168129929200650
 85901546975583577644230191713016530849872254943261333228208553748757868018711338287
 09313286601366479040704720680849935464542140720904086369896105835275219652626608375
 3869733406285692909824683

$\zeta_{\max} = 453$

$$\xi = 1, \Phi_1(908, 2^3) = 830833 \times 346 \ 637323 \times 20 \ 036217 \ 389373 \ 244862 \ 864196 \ 262187 \ 329665 \\ 281723 \ 794204 \ 087126 \ 504401 \ 910962 \ 446906 \ 740420 \ 971632 \ 225005 \ 619234 \ 545143 \ 765484 \ 571787 \\ 465453 \ 566822 \ 046032 \ 707843 \ 657710 \ 450847 \ 573608 \ 842023 \ 091477 \ 564503 \ 419913 \ 811228 \ 599412 \\ 238447 \ 477486 \ 057688 \ 078872 \ 719911 \ 924371 \ 752128 \ 031286 \ 077715 \ 556919 \ 687537 \ (260 \text{ digits})$$

T=0.9s

6.v=910

$\Pi_1 =$

23081540261671305056631293429525888216546118906068840049452022017848067251971680260
 34360618790233431057692076685206612339948901977304533291283421499503147207484535314
 83725314640546591616281888272339974185816856288361634547958442334110087861050643350
 15478933625142771639298731

$\zeta_{\max} = 455$

$\xi = 2, \Phi_1(910, 2^5) = 3^2 \times 14341 \times 4775591 \times 37446886656847465991689801239315679$
 624739 231358 928800 721195 922116 731834 107587 307131 483040 971907 188058 407841
 315162 734199 983990 866115 713353 711711 628902 689770 160555 698001 725152 373447
 996175 935434 554759 465933 006827 249223 042787 690762 164517 396758 215478 383294
 266960 578897 (263 digits)

T=0.4s

7.v=912

$\Pi_1 =$
 92326161046685220226525173718103552866184475624275360197808088071392269007886721041
 37442475160933724230768306740826449359795607909218133165133685998012588829938141259
 34901258562186366465127553089359896743267425153446538191833769336440351444202573400
 61915734500571086557194923

$$\xi_{\max} = 457$$

$\xi = 2, \Phi_2(912, 2^5) = \Phi_1^*(912, 2^5) = 23 \times 149 \times 659 \times 60719 \times 11041001 \times 26876151571 \times 2836188$
 753129 069940 492734 226989 316227 184045 585740 335285 079297 169982 072938 028004 816032
 845909 190630 974616 004415 530717 300703 618346 044670 780911 714583 808861 379615 020293
 776481 905988 243508 215935 481111 593654 563054 563999 425532 469293 817706 631274 780775
 366553 (247 digits)

T=1.0s

8.v=914

$\Pi_1 =$
 36930464418674088090610069487241421146473790249710144079123235228556907603154688416
 54976990064373489692307322696330579743918243163687253266053474399205035531975256503
 73960503424874546586051021235743958697306970061378615276733507734576140577681029360
 247662938002284346228779691

$$\xi_{\max} = 459$$

$\xi = 6, \Phi_2(914, 2^{13}) = \Phi_1^*(914, 2^{13}) = 29 \times 61 \times 1949 \times 25601 \times 5229957750865496264357355$
 480088 625174 001536 852714 167066 864397 704695 999698 801925 859661 277596 675508 722238
 218991 071382 073662 384074 750194 404457 201630 701080 016677 936353 199094 775345 829505
 129480 596437 475996 233880 466205 440690 156352 658022 840786 124621 665088 846591 316666
 865841 (265 digits)

T=0.2s

9.v=916

$\Pi_1 =$

14772185767469635236244027794896568458589516099884057631649294091422763041261875366
61990796025749395876922929078532231897567297265474901306421389759682014212790102601
49584201369949818634420408494297583478922788024551446110693403093830456231072411744
0990651752009137384915118763

$\xi_{\max} = 461$

$\xi = 1, \Phi_2(916, 2^3) = \Phi_1^*(916, 2^3) = 3 \times 5 \times 643 \times 1879 \times 15413 \times 5074187 \times 41543569 \times 171086$
 $649479 \times 183295 663130 862074 067638 363019 172742 346164 777991 554895 205082 555475$
 $491571 720023 959818 603684 236280 364495 760895 496783 666104 655742 353071 597443 639020$
 $055221 481430 590963 609534 893845 473115 725480 837024 043233 943621 581429 100015 236844$
 $781051 117821 631559$ (240 digits)

T=0.9s

10.v=918

$\Pi_1 =$

59088743069878540944976111179586273834358064399536230526597176365691052165047501466
47963184102997583507691716314128927590269189061899605225685559038728056851160410405
98336805479799274537681633977190333915691152098205784442773612375321824924289646976
3962607008036549539660475051

$\xi_{\max} = 463$

$\xi = 17, \Phi_1(918, 2^{35}) = 7^2 \times 11 \times 83 \times 132 080253 637656 840970 507882 020668 068565 970146$
 $410211 302784 266214 466081 883374 047212 523145 565954 422459 787372 682931 670181 898434$
 $610927 462717 777805 520241 382481 097327 948016 317217 695750 676011 922498 521129 532676$
 $289355 824507 277149 661974 118151 713754 816883 037972 236909 202853 042215 034987$ (273 digits)

T=0.1s

From the study above a method for calculating big prime numbers arises: we choose consecutive big even numbers v and apply the method for small values of ξ .

Theorem 4.1 highlights additional symmetries of the internal structure of the natural numbers. These symmetries determine a priori the signs of $\beta_i = \pm 1, i = 0, 1, 2, \dots, v-1$ in equation (4.1). In this article we will not make any more mention in other symmetries of equation (4.1).

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